

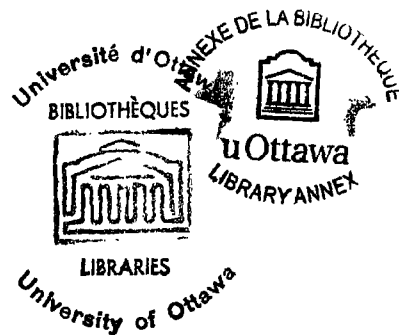
CUADRO: AN INTERACTIVE PACKAGE TO SOLVE THE QUADRATIC  
PROGRAMMING PROBLEM

by

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of the University of Ottawa  
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## ABSTRACT

The quadratic programming problem (QPP) deals with the optimization of a quadratic objective function, subject to linear constraints. In general, these constraints can be either in equality or inequality form and the nonnegativity conditions on the variables can apply for all or part of them. However, this problem has to be standardized as a preliminary step before trying to get a solution.

Some methods of solution are briefly presented and their computational behaviour analyzed. The method of Beale is widely discussed in both theoretical and practical ways and is adopted as the procedure to find the optimum. An interactive software package written in FORTRAN is presented.

## ACKNOWLEDGEMENTS

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## INTRODUCTION

Mathematical programming is the name given to the analysis of problems where the optimum of a function is to be found, and the variables are subject to inequality or equality constraints. The term "linear programming", corresponds to the case where the function to be optimized, i.e. the objective function, is linear and where the constraint set is the intersection of a finite number of half spaces. Nonlinear programming is, then, the case where the objective function is nonlinear.

An important case of nonlinear programming is the quadratic programming problem (QPP), where the objective function is a quadratic function and the constraints set is as in the linear case, a finite set of half spaces.

Research on nonlinear programming has gone almost parallel to the development of linear programming, but the nonlinear case has found some mathematical difficulties notably more important that have hindered its development. Thus, we still cannot speak today of a complete theory of nonlinear programming. In this development, the period of the last two decades has, thus far, appeared to be a time of appraisal and consolidation. Generally speaking, the major achievements in this time have been related to evaluating and improving the efficiency of existing solution methods.

The rates at which many of the most commonly used algorithms converge to optimal solutions have been subjected to both theoretical and empirical examination [8,25,29]

Recently, particular attention has been paid to "Large-scale Mathematical Programming", that is, to the difficulties encountered in attempting to solve the very large nonlinear problems that sometimes arise in Operations Research. In this field, important works have been done [16] in the developing of methods based on solving a sequence of QPP as a route towards the solution of a general nonlinear programming problem.

This thesis begins with a brief presentation of some of the most commonly used algorithms for solving the QPP and analyses their computational behaviour. The main intent of this thesis, however, is to present an interactive software package for the QPP, and discuss the most important aspects and problems encountered during its development.

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## Chapter I

### THE QUADRATIC PROGRAMMING PROBLEM

#### 1.1 HISTORICAL OUTLINE

The subject of mathematical programming has formally existed as a well defined subject area, since the late 1940's when George Dantzig developed the simplex method for linear programming. Later, in 1951 H.W.Kuhn and J.W.Tucker, in the Second Berkeley Symposium on Mathematical Statistics and Probability, published the paper "Nonlinear Programming" [18] which established necessary conditions for optimal solutions to nonlinear programming problems. Since that time, the history of Mathematical Programming can be divided, roughly speaking, into two parallel and occasionally convergent streams of development : one inspired by the Kuhn-Tucker optimality criterion and the other based on the principles of Dantzig's Simplex Method.

Initially, the development of Dantzig's Simplex Method was rather narrowly confined to the identification and study of various important subclasses of linear programming problems (Transportation, Network flows, Game theory, etc). Then in the late 1950's, a confluence of the Simplex

computational approach and the Kuhn-Tucker theory produced a number of different algorithms for solving quadratic programming problems. Basically, these methods rely on the fact that the first partial derivative of a quadratic function is linear : this is precisely what makes quadratic programs the most easily handled of all nonlinear problems.

Perhaps the most important period in the development of the solution methods of quadratic problems and, in general, of nonlinear ones was between 1959-1967. Some important works appeared presenting useful and efficient algorithms that can be applied to many nonlinear programs [1,11,30]. Some were applicable to problems having linear constraints only and still others were designed specifically for certain categories of the nonlinear programming problem... the quadratic programming problem is in this class.

## 1.2 STATEMENT OF THE PROBLEM

Generally speaking, the problem of optimizing a function of one or more variables is called a mathematical programming problem. In considering the nonlinear programming problem one always prefers to think of a minimization problem. This implies no loss of generality

since if the objective is to maximize  $F(x)$ , this can be converted to an equivalent minimization problem by letting  $f(x) = -F(x)$  and then minimizing  $f(x)$ .

Consider a convex quadratic function of  $n$  variables which is to be minimized :

$$G(x) = (1/2) x'Ax + b'x + c$$

where :

$x$  is an  $n$ -vector of the unknowns

$A$  is a symmetric positive semidefinite matrix with  $n$  rows and  $n$  columns

$b$  is an  $n$ -vector

$c$  is a scalar

The  $m$  constraints ( $m \leq n$ ) associated with the problem, are

$$Px - q \leq 0$$

$$x > 0$$

where :

$P$  is a matrix with  $m$  rows and  $n$  columns

$q$  is an  $m$ -vector with strictly positive components

Thus the QPP is characterized by linear constraints, nonnegativity requirement on the variables and an objective function  $G(x)$  which is a positive semidefinite quadratic form.

The reason why the matrix  $A$  must be positive semidefinite arises from one of the most serious problems in nonlinear programming: the problem of local and global minima. If

is a positive semidefinite matrix, any local minimum is also a global minimum. When A is not positive semidefinite, then there is a possibility that several local minima exist.

The QPP is often written as

$$\min \{ G(x) = (1/2)x'Ax + b'x + c \mid Px \leq q, x \geq 0 \} \quad (P1)$$

Note that the constraints and nonnegativity conditions are the same as in the case of linear programming. Constraints which are not in this form must be first transformed to this standard form by using any of techniques of linear programming; i.e. the phase I-phase II procedure or the big M technique (by adding artificial variables) [7].

In the statement of the QPP there are no restrictions on the choice of the elements of P, q or b; but the restriction that A be positive semidefinite is crucial. So far, there is no method for solving the QPP where A is an arbitrary symmetric matrix. As noted in the subsequent discussions, some methods even require that A be positive definite.

### 1.3 SOME APPLICATIONS

Although linear functions are the most widely used type in the formulation of mathematical optimization problems, quadratic functions are quite firmly established in second

place. This is due to the fact that a large number of the functional relationships occurring in the real world are either truly quadratic or may be so approximated.

For example, the area of a disk, cube or other regular figures is proportional to the square of their characteristic linear dimensions. The kinetic energy carried by a rocket or atomic particle is proportional to the square of its velocity, while the potential energy of a rigid standing wall or dam is a quadratic function of its height. The revenue of a monopolistic firm that sells  $x_1$  units of some product at a unit price of  $x_2$  is  $x_1 x_2$ , which is also quadratic. In statistics, the variance of a given sample of observations is a quadratic function of the values that constitute the sample.

One very common example of a QPP arises when data are to be fitted to a mathematical model by the least squares method. Suppose an economist theorizes that the fraction  $B$  of the American public's total consumption expenditure that is allocated to goods and services in the  $j^{\text{th}}$  class ( $j = 1, 2, \dots, n$ ) is constant from year to year, regardless of the overall level of consumption (assume all types of goods and services have been partitioned into  $n$  mutually exclusive classes). He has available as data a detailed breakdown of all consumption expenditures over an  $m$ -year period and wishes to use these statistics to form least squares estimates of

the unknown parameters  $B_j$ . Let  $C_i$  be the total consumption expenditure and let  $X_{ij}$  be the amount allocated on goods and services in the  $j^{\text{th}}$  class ( $j = 1, 2, \dots, n$ ) during the  $i^{\text{th}}$  year ( $i = 1, 2, \dots, m$ )

$$C_i = \sum_{j=1}^n X_{ij}$$

$j = 1, 2, \dots, n$   
 $i = 1, 2, \dots, m$

Then the problem can be formulated as

minimize  $Z$

where

$$Z = \sum_{i=1}^m \sum_{j=1}^n (X_{ij} - B_j C_i)^2$$

$i = 1, 2, \dots, m$   
 $j = 1, 2, \dots, n$

subject to  $\sum_{j=1}^n B_j = 1$  and  $0 \leq B_j < 1$

Another quadratic program known as the Portfolio problem arises in the planning of investments. Suppose one wants to invest a total of  $\$B$  in  $n$  different stocks and bonds  $x_j$  ( $j = 1, 2, \dots, n$ ). The annual return on a  $\$1$ -investment in the  $j^{\text{th}}$  security is a random variable whose expected value and variance (based on historical data) are  $F_j$  and  $S_{jj}$ ; the covariance of the  $i^{\text{th}}$  and  $j^{\text{th}}$  returns is  $S_{ij} = S_{ji}$ .

If one wants only to maximize the expected overall return, regardless of risk, it would only be necessary to invest in the security offering the greatest  $F_j$ . However, one might prefer to invest in such a way as to minimize the

overall variance of the portfolio subject to the condition of an expected total return of at least \$P per year. If we define S to be a symmetric N by N matrix with elements s the problem can be formulated as follows :

find the values of  $x_1, x_2, \dots, x_n$  that minimize Z where

$$Z = x'Sx$$

subject to

$$\sum_{j=1}^n R_j x_j \geq P$$

$$\sum_{j=1}^n x_j = B \quad \text{and} \quad x_j \geq 0$$

$$j = 1, 2, \dots, n$$

Chapter II  
MATHEMATICAL BACKGROUND

2.1 GLOBAL EXTREMA

A function  $f(x)$  takes its absolute minimum at a point  $x^*$  if  $f(x) \geq f(x^*)$  for all  $x$  over which the function  $f(x)$  is defined. The absolute minimum is also called global minimum.

A function  $f(x)$  has a relative minimum at some point  $x^0$  if there exists an interval including  $x^0$ , no matter how small, such that for all  $x$  in this interval,  $f(x) \geq f(x^0)$ . The relative minimum is also referred to as local minimum. Similarly, there exist global maxima and local maxima, whose definitions follow logically from the previous ones.

For a concave function, linear interpolation between any two points never overestimates the value of the function. Obviously, if  $f(x)$  is concave, then  $-f(x)$  is convex.

Let  $f(x)$  be a convex function over a closed interval  $a \leq x \leq b$ , then any local minimum of  $f(x)$  in this interval is also the global minimum of  $f(x)$  over the interval. The global maximum of a convex function  $f(x)$  over a closed

interval  $a \leq x \leq b$  will be taken on at either  $x = a$  or  $x = b$  or both. The two previous statements are often stated as theorems and their proofs can be found in the optimization literature [9,10]. We may well be interested in minima that occur when one or more of the variables are at their upper or lower bound. It is readily seen that there are  $2^n$  possible combinations of the  $a_i$  and  $b_i$  which the  $x_i$  may take. If it is to minimize  $z = f(x)$  and there is no characteristic known of  $f(x)$  (even if it is known that it is continuous and differentiable), it can only be hoped to find a local minimum. However, considerably more can be said if the function being maximized or minimized is concave or convex.

If  $f(x)$  is a convex function, it has a unique minimum over a feasible convex region and the same minimum value is found for a convex subset of the feasible region. Furthermore, the minimum of a concave function over a convex region occurs at an extreme point of that region, or otherwise points which are not extreme points can be written as linear combinations of extremal points.

The problem (P1) is said to be a convex quadratic programming problem if  $A$  is positive semidefinite, so that  $G(x)$  is convex. If  $A$  is negative semidefinite, so that  $G(x)$  is concave, then problem (P1) is called a concave quadratic programming problem.

## 2.2 THE SIMPLEX METHOD

### 2.2.1 The canonical form

When dealing with linear programming problems, the simplex method provides the most efficient procedure for solution. To apply the simplex method, the problem is assumed in its canonical form

$$\begin{aligned} \text{Max } z &= b'x \\ Px &= q && \text{(P2)} \\ x &\geq 0 \end{aligned}$$

with  $n$  unknowns and  $m$  constraints, then :

$b$  is an  $n$ -vector,

$P$  is an  $m \times n$  matrix,

$q$  is an  $m$ -vector and

$x$  is the  $n$ -vector of the unknowns.

This formulation is clearly very restrictive; however linear programming applies equally to situations where the constraints are inequalities or a mixture of equalities and inequalities, where some of the variables can have negative values and where the objective function is to be maximized.

In general, the case may be

$$\text{maximize } Z$$

where

$$Z = d'y$$

Subject to

$$Fy \leq r$$

$$Gy \geq s$$

$$Hy = t$$

$$y_i \geq 0 \quad i = 1, 2, \dots, k \quad ; \quad k \leq n$$

which need to be transformed into canonical form.

If a constraint is a less-than condition, then an equation can be made out of this constraint by adding a nonnegative variable  $u$  to the left-hand side and writing :

$$Fy + u = r$$

$u$  is called a slack variable and is an additional unknown that has to be determined. Similarly, if a constraint is a greater-than condition, it can be written as

$$Gy - v = s$$

where  $v$  again is a nonnegative variable (usually called a surplus variable) that constitutes an additional unknown of the problem. If some of the variables are not constrained to be nonnegative, then these variables can be expressed as

$$y_j = y_j^+ - y_j^-$$

$$y_j^+ \geq 0 \quad y_j^- \geq 0$$

$$j = 1, 2, \dots, n-k$$

because any number can always be expressed as the difference of two nonnegative numbers. Finally, instead of maximizing  $Z$ ,  $-Z$  will be minimized.

The problem becomes

minimize  $Z$

where

$$z = \begin{bmatrix} d & 0 & 0 \end{bmatrix} \begin{bmatrix} y \\ u \\ v \end{bmatrix}$$

subject to

$$\begin{bmatrix} F & I & 0 \\ G & 0 & -I \\ H & 0 & 0 \end{bmatrix} \begin{bmatrix} y \\ u \\ v \end{bmatrix} = \begin{bmatrix} r \\ s \\ t \end{bmatrix}$$

By putting

$$x = \begin{bmatrix} y \\ u \\ v \end{bmatrix} \quad P = \begin{bmatrix} F & I & 0 \\ G & 0 & -I \\ H & 0 & 0 \end{bmatrix} \quad q = \begin{bmatrix} r \\ s \\ t \end{bmatrix} \quad b = \begin{bmatrix} -d \\ 0 \\ 0 \end{bmatrix}$$

problem (P2) is gotten, which is in canonical form.

It should be noted that the number of unknowns have been increased by  $u + v + (n-k)$ . Thus any method that solves the canonical form will also handle more general situations involving mixtures of equalities and inequalities, partial nonnegativity constraints, and minimization or maximization. The canonical form of the problem will be referred to hereafter.

### 2.2.2 Basic solutions

It will be convenient to define some concepts that can be useful in discussing future topics.

1. Feasible solution : Any solution to  $Px = q$  ,  $x \geq 0$
2. Basic solution : A solution of  $Px = q$  , obtained by setting  $n-m$  variables equal to zero and solving the remaining  $m$  variables, provided that the determinant of the coefficients of these  $m$  variables does not vanish (so that the values of these  $m$  variables are uniquely determined).
3. Basis : The collection of  $m$  variables which are not equal to zero in the construction of a basic solution.
4. Basic feasible solution : A basic solution of  $Px = q$  which also satisfies  $x \geq 0$ .
5. Optimal solution : A feasible solution which satisfies  $b'x = \min$ .
6. Optimal basic solution : A basic feasible solution which satisfies  $b'x = \min$ .

Before starting the simplex procedure, a basic feasible solution must be found. For this, the vector  $q$  in the constraints  $Px = q$  must be strictly positive. Whenever necessary, a new transformation has to be done to the system such that one decision variable be isolated in each constraint with a +1 coefficient, that that variable not

appear in any other constraint and that it have a zero coefficient in the objective function.

The variables added at this point are called artificial variables because they don't belong to the system and must eventually be suppressed. They are introduced only because they constitute immediately a basic feasible solution. After creating the artificial variables, they have to vanish. For this, the infeasibility form  $w$  defined by

$$x_{n+1} + x_{n+2} + \dots + x_{n+m} = w$$

has to be driven to zero [7,13]. Only then have we a basic feasible solution to start the simplex procedure.

### 2.2.3 Computational criteria

1. Optimality criterion : Suppose that in a minimization problem every nonbasic variable has a nonnegative coefficient in the objective function of a canonical form. Then the basic feasible solution given by that canonical form minimizes the objective function over the feasible region.
2. Unboundedness criterion : Suppose that in a minimization problem some nonbasic variable has a negative coefficient in the objective function of a canonical form. If that variable has positive or zero coefficients in all constraints, then the objective

function is unbounded from below over the feasible region.

3. Improvement criterion : Suppose that in a minimization problem some nonbasic variable has a negative coefficient in the objective function of a canonical form. If that variable has a negative coefficient in some constraint, then a new basic feasible solution may be obtained by pivoting.
4. Ratio and pivoting criterion : When improving a given canonical form by introducing the variable  $x_5$  into the basis, pivot in a constraint that gives the minimum ratio of the right-hand-side coefficient to the corresponding  $x_5$  coefficient. Compute these ratios only for constraints that have a positive coefficient of  $x_5$ .

When introducing the variable  $x_5$  into the basis, another basic variable must leave its place to  $x_5$ . This substitution (pivoting)<sup>1</sup> is merely the familiar variable elimination technique from high-school algebra, known more formally as Gauss-Jordan elimination. Consequently, after pivoting, the form of the problem has been altered, but the modified equations still represent the original problem and have the same feasible solutions and the same objective value when

---

<sup>1</sup> The commonly known procedure for pivoting is not treated in this paper. However, it can be found in any book dealing with linear programming [7,13].

evaluated at any given feasible point.

#### 2.2.4 Geometrical interpretation

Consider the following problem stated in its canonical form :

minimize  $Z$

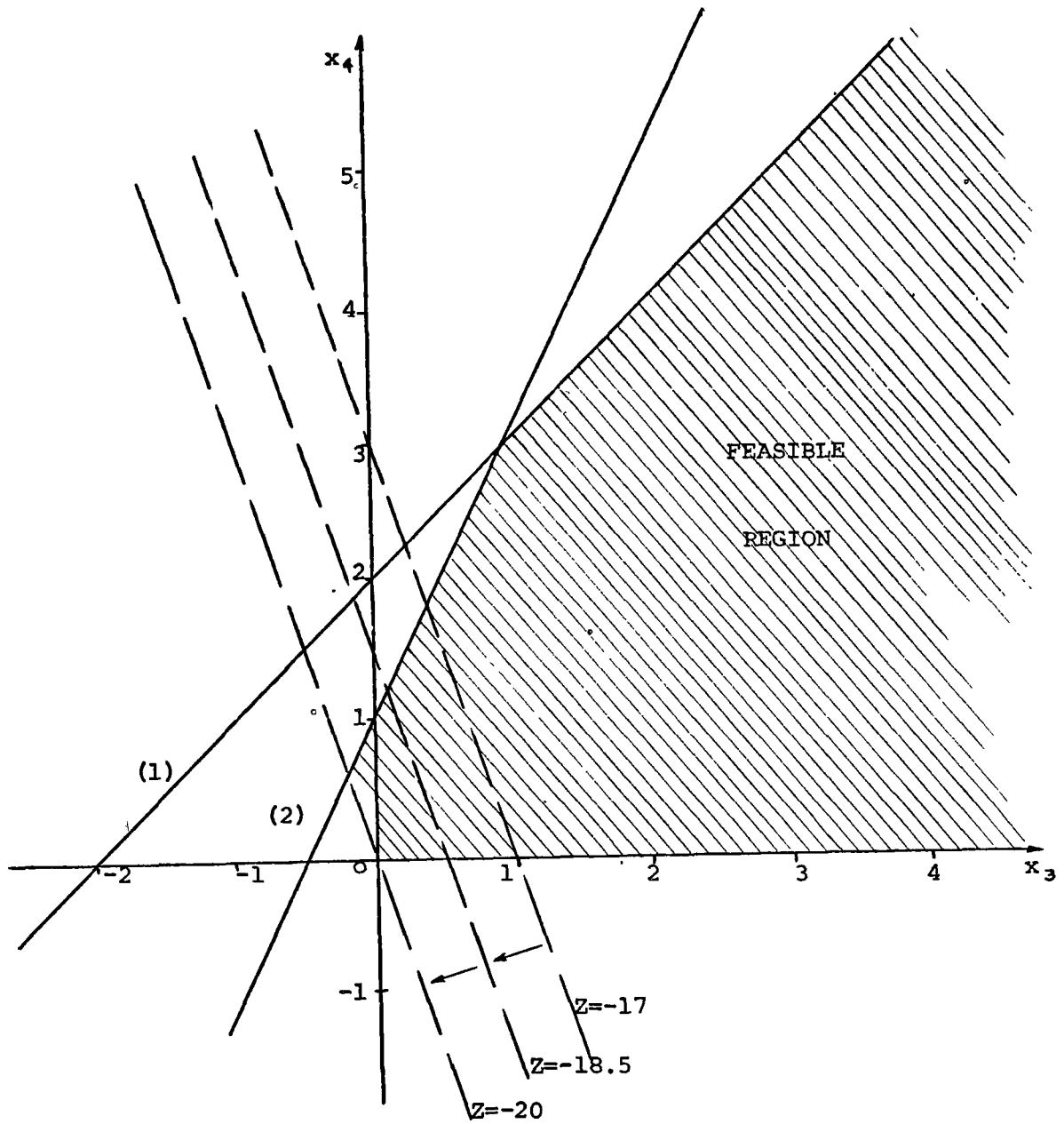
where

$$\begin{aligned} Z &= 0x_1 + 0x_2 + 3x_3 + x_4 - 20 \\ \text{subject to } x_1 & - 3x_3 + 5x_4 = 6 \quad (1) \\ x_2 & - 8x_3 + 4x_4 = 4 \quad (2) \\ x_j &\geq 0 \quad (j = 1, 2, 3, 4) \end{aligned}$$

Figure 1 illustrates the situation. For any value of  $z$ , say  $z = -17$ , the objective function is represented by a straight line. As  $z$  decreases to  $-20$ , the line corresponding to the objective function moves parallel to itself across the feasible region. At  $z = -20$  it meets the feasible region only at the point  $x_1 = x_2 = 0$  and, for  $z < -20$  it no longer touches the feasible region. Consequently,  $z = -20$  is optimal.

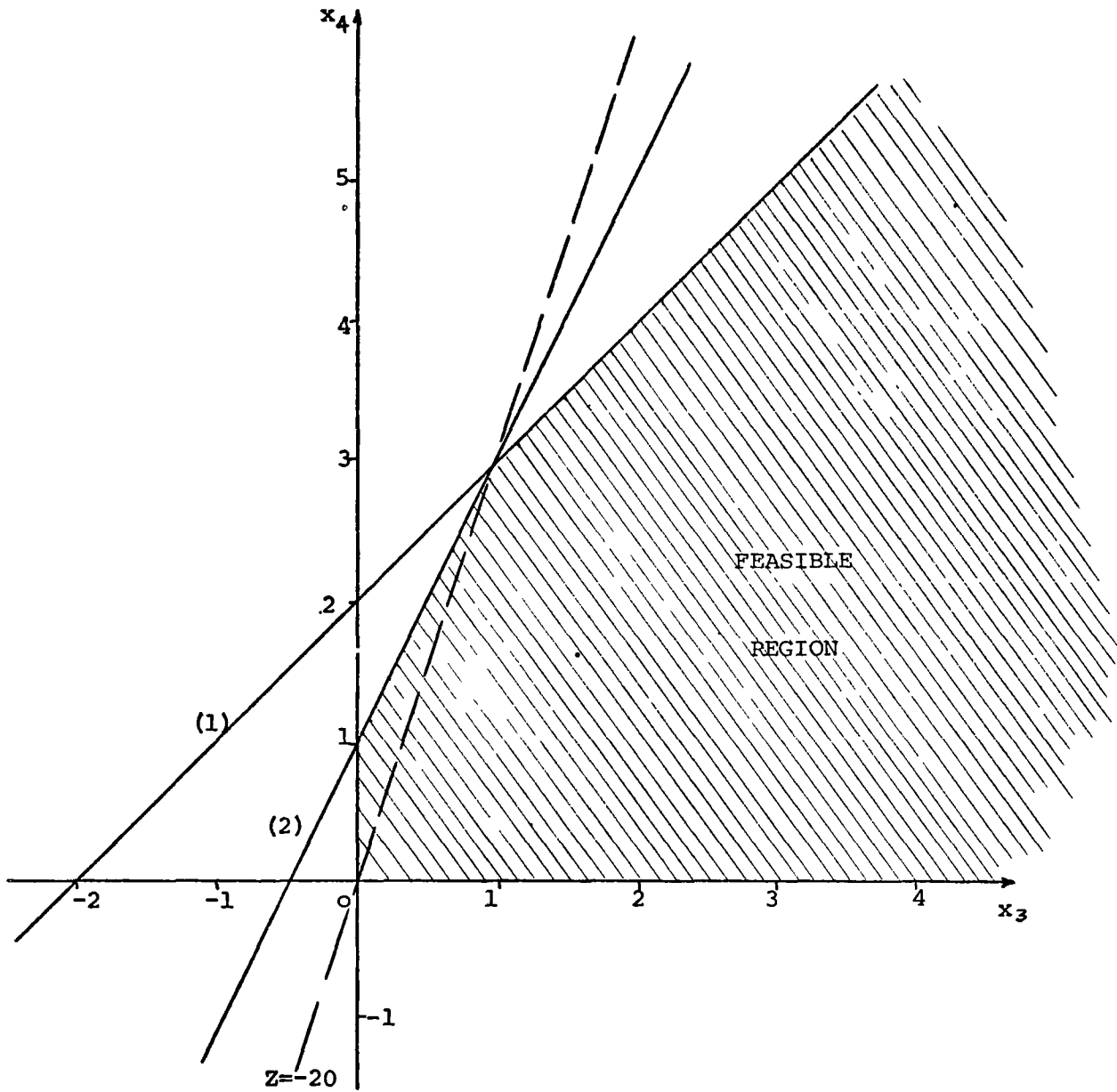
If we change our objective function to  $z = -3x_3 + x_4 - 20$  (which is shown in figure 2) decreasing  $x_3$ , while holding  $x_4 = 0$ , corresponds to moving outwards from the origin along the  $x_3$ -axis. As we move along the axis, we never meet either constraint (1) or (2). Also, as we move along the  $x_3$ -axis, the value of the objective function is increasing

Figure 1: GEOMETRICAL INTERPRETATION OF EXAMPLE 1



to  $+\infty$ . Hence the objective function is unbounded from

Figure 2: GEOMETRICAL INTERPRETATION OF EXAMPLE 2



above.

On the other hand, if our objective function is changed to  $z = 3x_3 - x_4 - 20$ , which is also shown in figure 2, decreasing  $x_4$  corresponds to moving from the origin along the  $x_4$ -axis. In this case, however, we encounter constraint 2 at  $x_4 = 1$  and constraint 1 at  $x_4 = 2$ . Consequently, to maintain feasibility in accordance with the ratio criterion, we move to the intersection of the  $x_4$ -axis and constraint 2, which is the optimal solution.

### 2.3 THE KUHN-TUCKER CONDITIONS

In attempting to develop algorithms for solving nonlinear programming problems, it is useful to have some information concerning the characteristics of an optimal solution. Kuhn and Tucker [18] showed that if a vector  $x^*$  minimizes  $f(x)$ , subject to constraints  $h_i(x) \geq 0$ , then,

$$\begin{aligned} \frac{\partial f}{\partial x_j} + \sum_{i=1}^m t_i \frac{\partial h_i}{\partial x_j} &= 0 \\ h_i(x^*) &\geq 0 && \text{(KT1)} \\ t_i h_i(x^*) &= 0 \\ t_i &\geq 0 \end{aligned}$$

for all  $i = 1, 2, \dots, m$   
 $j = 1, 2, \dots, n$

where

$t_i$  is the Lagrange multiplier associated with the  $i^{\text{th}}$  constraint.

These are the well-known Kuhn-Tucker conditions.

However, in the case of the QPP, the Kuhn-Tucker conditions may have a special form. The function  $f(x)$  to be minimized is in this case

$$G(x) = (1/2)x'Ax + b'x + c$$

and the constraints  $h(x) \geq 0$  are of two kinds

The side constraints

$$h^1 = q - Px \geq 0$$

and the nonnegativity constraints

$$h^2 = x \geq 0.$$

Then,

$$\partial G(x) / \partial x = Ax + b$$

$$\partial h^1(x) / \partial x = \partial / \partial x (q - Px) = - \sum_{i=1}^n u_i P_{ij} = -P'u$$

where  $u_i$  is the Lagrange multiplier associated with the  $i^{\text{th}}$  side constraint. Clearly, the  $m$ -vector  $u$  is the Lagrange multiplier associated with the  $m$  side constraints  $h^1(x)$ .

Similarly, for the nonnegativity constraints we have

$$\partial h^2(x) / \partial x = \sum_{j=1}^n v_j e_j = v$$

where  $e_j$  is the  $j^{\text{th}}$  column of the identity matrix  $I$  and  $v_j$  is the Lagrange multiplier associated with nonnegativity constraint  $x_j$ .

At the minimum point, the conditions are :

$$Ax + b = -P'u + v$$

$$u \geq 0$$

$$v \geq 0$$

$$x \geq 0$$

Moreover,  $t_i h_i(x^*) = 0$  becomes

$$u_i h_i(x^*) = 0 \quad \text{or} \quad u_i (q_j - \sum_{i=1}^n P_{ij} x_i) = 0$$

$j = 1, 2, \dots, n$

If  $y_i$  represents the nonnegative slack variable associated with the  $i^{\text{th}}$  constraint

$$y_i = q_i - \sum_{j=1}^n (P_{ij} x_j)$$

then, from above

$$u'y = 0$$

similarly, for the nonnegativity constraints  $v'x = 0$ .

Then,  $t_i h_i(x^*) = 0$  becomes  $v'x + u'y = 0$  which is true.

Since  $v, u, x, y \geq 0$ , the only way for it to hold is that  $v'x = 0$  and  $u'y = 0$ .

Then, the Kuhn-Tucker conditions for the QPD can be written as

$$Ax + b + P'u = v$$

$$v'x + u'y = 0 \quad \text{(KT2)}$$

$$q - Px = y$$

$$x, y, u, v \geq 0$$

Chapter III  
THE METHOD OF BEALE

3.1 THEORETICAL ASPECTS OF THE METHOD

The method of Beale [1] was published as an application of the simplex method to quadratic programming and it reduces to the simplex method when the objective function is linear.

Consider the problem of minimizing the quadratic objective function  $G(\bar{x})$  ( $\bar{x}$  an  $n$ -vector), subject to the constraints

$$\begin{aligned} P\bar{x} &\leq q \\ \bar{x} &\geq 0 \end{aligned}$$

where  $P$  is an  $m \times n$  matrix and  $q$  is an  $m$ -vector.

We can assume without loss of generality that a basic feasible solution can be obtained from this set of constraints (see section 2.3.1). Call  $y$  the vector of the variables that give a basic feasible solution at the starting point, then the constraints become:

$$\begin{aligned} P\bar{x} + y &= q \\ \bar{x} \geq 0, y &\geq 0 \end{aligned}$$

or

$$\left[ \begin{array}{c|c} P & I \end{array} \right] \begin{bmatrix} x \\ y \end{bmatrix} = [q]$$

Let

$$[x] = \begin{bmatrix} \bar{x} \\ y \end{bmatrix} \quad \text{and} \quad [P] = [ \bar{P} | I ]$$

Then  $x$  is an  $(n+m)$ -vector and  $P$  an  $m \times (n+m)$ . As a result the constraints can be written

$$Px = q$$

$$x \geq 0$$

It is assumed that the basic variables are the last  $x_{n+k}$  ( $k=1, 2, \dots, m$ ) and the first  $x_j$  ( $j=1, 2, \dots, n$ ) are nonbasic. Define the vector  $z = (z_0, z_1, \dots, z_n)$

$$z = \begin{bmatrix} 1 \\ x \end{bmatrix}$$

Then the objective function can be expressed in terms of  $z$  as

$$G(z) = z' \Gamma z$$

where

$$\Gamma = \begin{bmatrix} c & | & b'/2 \\ \hline & & \\ \hline b/2 & | & (1/2)' \end{bmatrix}$$

Note now that for  $1 \leq k \leq m$  the  $(k+1)$ th entry of  $\partial G / \partial z_k$  is

$$b_k + \sum_{j=1}^n a_{kj} z_j$$

The constraints can be used to express the basic variables in terms of the nonbasic ones

$$x_i = p_{i0} - \sum_{j=1}^n p_{ij} z_j$$

where  $p_{i_0} = q_i$  at the starting point where a basic feasible solution is immediately available. At any trial solution of the problem, the basic variable  $x_i = p_{i_0}$  and the nonbasic variables  $z$  are all zero. Also, at any trial point, the equations of the constraints can be used to express the objective function  $G$  in terms of the nonbasic variables only.

If  $\partial G / \partial z_s \geq 0$ , then an increase in  $z_s$  will not reduce  $G$ ; but if  $\partial G / \partial z_s < 0$ , then a small increase in  $z_s$  will reduce  $G$ . In the general case, when  $G$  is an arbitrary continuously differentiable function, the process of increasing  $z_s$  (when  $\partial G / \partial z_s < 0$ ), must terminate when either one of the following conditions holds:

1. A basic variable, say  $x_t$ , has become zero and is about to become negative.
2.  $\partial G / \partial z_s$  vanishes and is about to become positive.

In case 1. (which is the only possibility in the case of linear programming), the remedy is to change the basis by making  $x_t$  nonbasic in place of  $z_s$ . Case 2. can give rise to substantial difficulty if  $G$  is an arbitrary function; but if  $G$  is a quadratic function, then  $\partial G / \partial z_s$  is a linear function of the nonbasic variables. In this situation, case 2. results in the creation of a new constraint.

If  $\partial G / \partial z_s$  becomes positive as  $z_s$  is increased, (with all other nonbasic variables equal to zero) then a new variable is defined as follows:

$$u = h_s + \sum_{j=1}^n a_{sj} z_j$$

Now  $u$  is made the new nonbasic variable, and throughout the constraints and the expression for  $G$ , the above equation is used to substitute for  $z_s$  in terms of  $u$  and the other nonbasic variables. Note that if  $z_s$  is one of the original  $x$ -variables (not one of the introduced  $u$ -variables), then there will be additional basic variable after this step.

One important feature of the new  $u$  variable is that it is not restricted to nonnegative values. It is therefore called a 'free' variable as opposed to the original  $x$ -variables which are called restricted variables. Since  $u$  is not sign restricted, if  $\partial G / \partial u > 0$  then  $G$  can be reduced by making  $u$  negative.

It is convenient to have a set of nonbasic variables that all vanish at the trial solution, since the values of the basic variables are then simply given as the constant terms in the tableau. The new variable is chosen, so that it will not be profitable to change its value when we change the values of any other nonbasic variable. It has been shown [19,3] that this method terminates in a finite number of steps.

### 3.2 UPDATING THE TABLEAUX

When passing from one iteration to the next, the expressions for the constraints are handled exactly as in linear programming, (the constraints part is often called the simplex tableau). However, for the objective function, there is a different approach.

TABLE 1  
INITIAL TABLEAU FOR BEALE

	variables	value	z
Constraints tableau	x	a	-P
objective function tableau		c	b/2
	z	b/2	$\frac{1}{2}P$

Table 1 shows the initial tableau for the method of Beale. The constraints tableau shows the basic variables x in terms of the nonbasic ones. The objective function tableau shows the symmetric matrix  $\Gamma$  of size  $(n+1) \times (n+1)$  whose first row can be used to determine optimality in the trial point: a positive value in each one of the columns v, where  $z_v$  is a restricted variable and a zero value in each

of the columns  $t$ , where  $z_t$  is a free variable (""), denotes an optimal tableau.

The procedure suggested by Beale can be expressed algebraically by saying that, starting from the objective function that is in terms of the nonbasic variables  $z$ , the  $z$ -variable that is to become basic is replaced by the equivalent expression in terms of the basic variable  $x$  (that is to become nonbasic) and the other nonbasic variables, giving rise to a new expression for  $G$  in terms of the new nonbasic variables. This is not very convenient in a computer, since it involves operating on both the rows and columns of the matrix. It is worthwhile examining the algebraic expressions for the combined effects of these transformations.

If we denote the pivotal column by  $s$  and if the expression for the new basic variable  $z$  is, in terms of the new nonbasic variables

$$z_t = e_0 + e_s z_s + \sum_{k \neq s} e_k z_k$$

where  $z_s$  denotes the new nonbasic variable replacing  $z_t$ , then the new coefficients  $g'_{kj}$  are given in terms of the old coefficients  $g_{kj}$  and the  $e_k$  as follows :

$$g'_{ss} = g_{ss} e_s^2$$

$$g'_{ks} = g'_{sk} = g_{ks} e_s + g_{ss} e_s e_k$$

$$g'_{kj} = g_{kj} + g_{ks} e_j + g_{sj} e_k + g_{ss} e_k e_j$$

where  $k, j \neq s$

If  $g^*$  is defined as

$$g^* = (1/2)g \quad e$$

$$g^* = g + (1/2)g \quad e \quad \text{for } k \neq q$$

then, the above expressions can also be written

$$g' = 2g^* \quad e$$

$$g' = g' = g^* \quad e + g \quad e$$

$$g' = g + g^* \quad e + g^* \quad e$$

The optimum point is reached when the derivatives of the objective function with respect to each of the nonbasic variables are nonnegative and the derivative of the objective function with respect to the free variables introduced in the problem are all zero. This is so because at this point an admissible change in any of the restricted variables will increase the value of the objective function and it is impossible either to increase or to decrease any of the free variables. Hence the value of  $G$  cannot be minimized further.

### 3.3 A NUMERICAL EXAMPLE

To illustrate this method, an example given by Beale [ 2 ] that has the characteristic of highlighting the principal features of the algorithm, is presented.

The problem is

$$\text{minimize } G = -6x_1 + 2x_1^2 - 2x_1x_2 + 2x_2^2$$

$$\text{subject to } x_1 + x_2 \leq 2 \quad x_1, x_2 \geq 0$$

Introducing a slack variable  $x_3$ , we have the constraint:

$$x_1 + x_2 + x_3 = 2 \quad x_1, x_2, x_3 \geq 0$$

A feasible solution for starting the process is

$$x_1 = 0 \quad x_2 = 0$$

This corresponds to point A in figure 3. At this point, the following expressions for the variable  $x_3$  and the objective function hold:

$$\begin{aligned} x_3 &= 2 - x_1 - x_2 \\ G &= (-3x_1) + (-3 + 2x_1 - x_2)x_1 \\ &\quad + (-x_1 + 2x_2)x_2 \end{aligned}$$

and so at the trial point:

$$x_1 = 0, \quad x_2 = 0, \quad x_3 = 2, \quad G = 0$$

Since  $\partial G / \partial x_1 = -3$ ,  $G$  can be decreased by increasing  $x_1$  ( $x_2$  held at its value of 0)

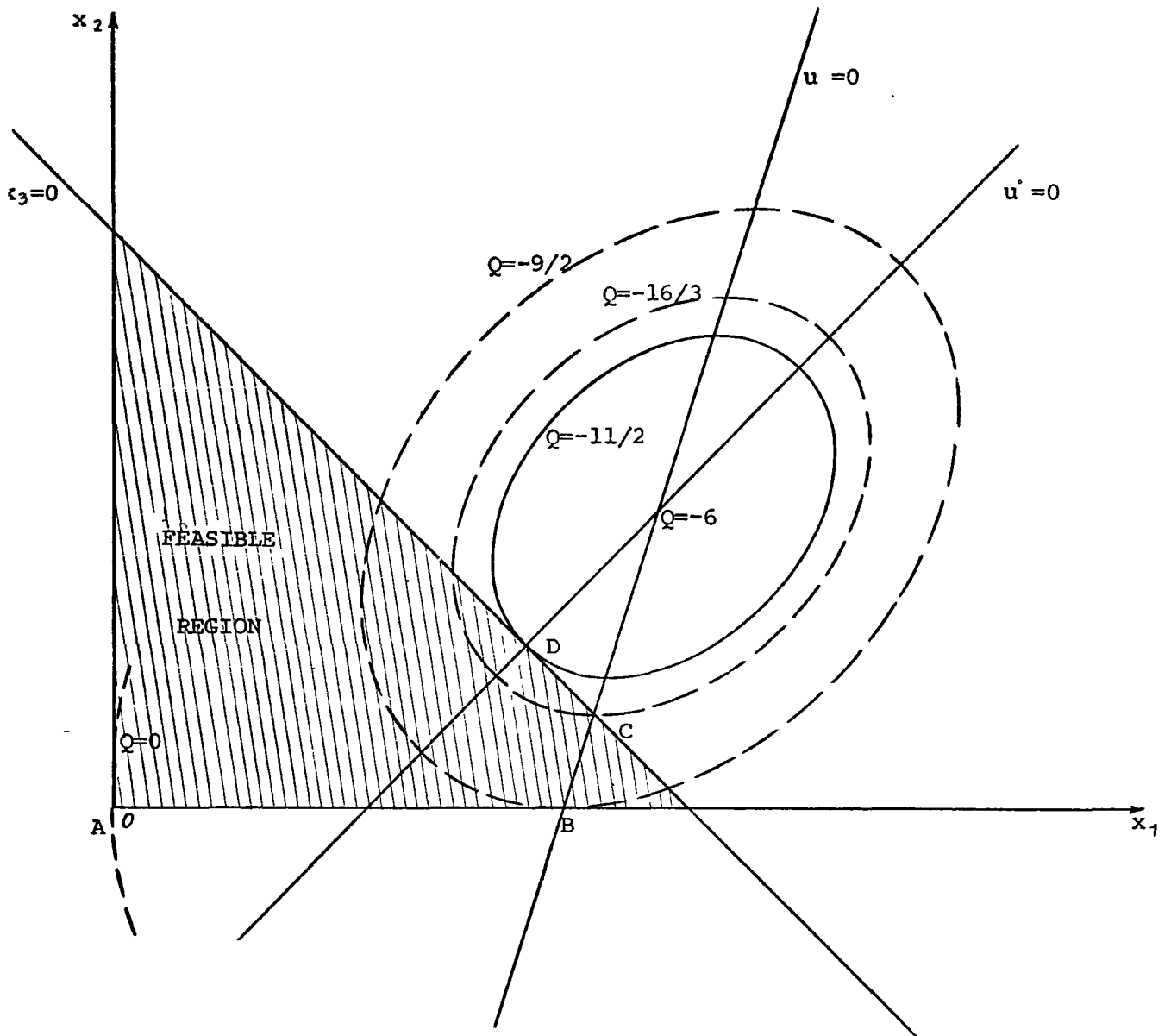
$$\partial G / \partial x_1 = -3 + 2x_1 - x_2 \quad \text{vanishes for } x_1 = 3/2$$

This point still belongs to the feasible region since  $x_3$  equals 0 only when  $x_1 = 2$  and then  $\partial G / \partial x_1 = -3 + 2x_1 = +1$  so, a free variable is introduced:

$$u_1 = \partial G / \partial x_1 = -3 + 2x_1 - x_2$$

where the subscript 1 on  $u$  means that it is the first free variable introduced into the problem. This  $u_1$  variable replaces  $x_1$  as an independent variable that vanishes.

Figure 3: THE METHOD OF BEALE ON EXAMPLE F



$u_1$  is zero in all the points where the ellipse  $G =$

constant has a horizontal tangent. From above :  $x_1 = 3/2 + (1/2)u_1 + (1/2)x_2$

Hence

$$s = 1 \quad e_0 = 3/2 \quad e_1 = 1/2 \quad e_2 = 1/2$$

$$(1/2)q_{s_5} = 1 \quad q_0^* = -3/2 \quad q_1^* = 1/2 \quad q_2^* = -1/2$$

Then the second trial point can be set and the tableau reads

$$\begin{aligned} x_1 &= 3/2 + (1/2)u_1 + (1/2)x_2 \\ x_2 &= 1/2 - (1/2)u_1 - (3/2)x_2 \\ G &= (-9/2 - (3/2)x_2) 1 \\ &\quad + ( (1/2)u_1 ) u_1 \\ &\quad + (-3/2 (3/2)x_2 ) x_2 \end{aligned}$$

At the second point B

$$u_1 = 0, \quad x_2 = 0, \quad x_1 = 3/2, \quad x_3 = 1/2, \quad G = -9/2$$

Since  $\partial G / \partial x_2 = -3/2$ ,  $x_2$  must be put into the basis ( $u_1$  stays 0) ~

$$\partial G / \partial x_2 = -3/2 + (3/2)x_2 \quad \text{vanishes for } x_2 = 1$$

It is not possible to reach this value since  $x_3$  vanishes before  $x_2$ , that is, for  $x_2 = 1/3$ .

Then we make  $x_3$  nonbasic and we express the new basic variable  $x_2$  and the previous one  $x_1$ , as functions of  $x_3$  and  $u_1$ , the new nonbasic set.

This gives as a third point :

$$\begin{aligned} x_2 &= 1/3 - (1/3)u_1 - (2/3)x_3 \\ x_1 &= 5/3 + (1/3)u_1 - (1/3)x_3 \end{aligned}$$

Hence,

$$s = 2 \quad e_0 = 1/3 \quad e_1 = -1/3 \quad e_2 = -2/3$$

$$(1/2)q_{ss} = 3/4 \quad q_0^* = -5/4 \quad q_1^* = -1/4 \quad q_2^* = -1/2$$

And the objective function at this point is :

$$\begin{aligned} G &= (-16/3 + (1/3)u_1 + (2/3)x_3) 1 \\ &\quad + (1/3 + (2/3)u_1 + (1/3)x_3) u_1 \\ &\quad + (2/3 + (1/3)u_1 + (2/3)x_3) x_3 \end{aligned}$$

The trial point corresponds to point C

$$u_1 = 0, \quad x_3 = 0, \quad x_2 = 1/3, \quad x_1 = 5/3, \quad G = -16/3$$

Since an x-variable has been made nonbasic,  $u_1$  must become basic. The only way to decrease  $G$  at this point is by decreasing  $u_1$ , because  $\partial G / \partial u_1 > 0$ . This is possible until

$$5/3 + (1/3)u_1 - (1/3)x_3 = 0$$

that is, for  $u_1 = -5$ . But  $\partial G / \partial u_1 = 1/3 + (2/3)(-5) = -3$  vanishes first, so the next nonbasic variable is a free variable  $u_2$  defined by

$$u_2 = 1/3 + (2/3)u_1 + (1/3)x_3$$

i.e.

$$u_1 = -1/2 + (3/2)u_2 - (1/2)x_3$$

hence,

$$s = 1 \quad e_0 = -1/2 \quad e_1 = 3/2 \quad e_2 = -1/2$$

$$(1/2)q_{ss} = 1/3 \quad q_0^* = 1/6 \quad q_1^* = 1/2 \quad q_2^* = 1/6$$

and a new tableau is available:

$$\begin{aligned}
u_1 &= -1/2 + (3/2)u_2 - (1/2)x_3 \\
x_1 &= 3/2 + (1/2)u_2 - (1/2)x_3 \\
x_2 &= 1/2 - (1/2)u_2 - (1/2)x_3 \\
G &= (-11/2 + (1/2)x_3) + 1 \\
&\quad + (3/2)u_2 + (1/2)x_3 \\
&\quad + (1/2)x_3
\end{aligned}$$

Since the derivative of G with respect to the free variable  $u_2$  is zero and the derivative with respect to the restricted variable  $x_3$  is positive, the minimum point has been reached. The problem solution therefore is:

$$x_1 = 3/2, \quad x_2 = 1/2, \quad x_3 = 0, \quad G = -11/2$$

This corresponds to point D in figure 3

### 3.4 THE ALGORITHM

1. Optimality : Check the first row of the objective function tableau. If the tableau is optimum, stop.
2. Pivot column : Select in the first row of the tableau, among the u-column a non-zero element. If all the elements in the first row corresponding the u-columns are zero or there are no u-columns then choose among the x columns the most negative. Call

this element  $g$ . Then  $r$  is the pivot column, which defines the variable  $z$  (either  $u$  or  $x$ ) that becomes basic.

3. Pivot row : Divide the absolute value of the element  $g$  by the element  $a$  in the diagonal of the objective function tableau. If this element is positive, then divide the elements  $p$  of the first column of the constraints tableau by the absolute value of the correspondent element  $p$  in the pivot column, but only if this element has the same sign as the element  $g$ . Obtain the same ratio in all the rows (except the first one) of the objective function tableau having in the first column one element with the same sign as  $g$ , (take absolute value of  $g$ ,  $h \neq 0$ ). The row giving the minimum ratio is selected as pivot row. The element located in the intersection of the pivot column and the pivot row (either  $g$  or  $p$ ) is called the pivot element.
4. If the pivot element is  $p$ , that is the pivot row belongs to the constraints tableau, replace the variable  $z$  (in the top of the pivot column) by the variable  $z$  (heading the pivot row).
5. If the pivot element is  $g$ , that is the pivot row belongs to the objective function tableau, replace  $z$

by a new variable  $u$  , that goes to the constraints tableau.

6. Update the tableaux (constraints section and objective function section) and repeat from step 1.

## Chapter IV

### SOME METHODS FOR THE QPP AND THEIR COMPUTATIONAL PERFORMANCE

Several articles have appeared in the literature comparing the computational efficiency of three of the best known simplicial methods of quadratic programming [1,11,30], namely the methods of Wolfe, Beale and Dantzig. A large number of problems have been solved with coefficients which were random variates obtained from a pseudo-random number generator. In certain cases, problem parameters were varied to determine whether they have a significant effect on the rate of convergence of the algorithms.

The following results are taken from a study carried out by Braitsch [8], who followed an experimental approach to compare the algorithms of Wolfe, Beale and Dantzig. A fourth method is also treated which is a slight modification of Wolfe's algorithm. This modification have been introduced in an attempt to reduce the effects of some of the inefficiencies inherent in Wolfe's algorithm. The method of Beale was presented in the previous chapter. A brief discussion of the methods of Wolfe and Dantzig is presented here.

#### 4.1 THE ALGORITHM OF WOLFE

Wolfe [30], working from an idea suggested by Makowitz [23], developed a method for the QP that is based on solving the Kuhn-Tucker conditions. Computationally, it corresponds to a simple modification of the artificial variable technique for finding a first feasible solution to a linear programming problem [7].

For the problem (P1),  $(3n + 2m)$  supplementary variables are added to the corresponding Kuhn-Tucker conditions (KT2)

$$w_i \quad i = 1, 2, \dots, n$$

$$z_j^1, z_j^2 \quad j = 1, 2, \dots, (m+n)$$

Now note that the  $x$ -vector which satisfies the following conditions when  $w = z^1 = z^2 = 0$  and  $r = 1$ , is the solution to the original problem:

$$Px + y \quad + \quad w = q$$

$$Ax \quad - \quad v + P'u + z^1 - z^2 + rb = 0$$

$$v'x + u'y = 0$$

$$x, v, y, u, w, z^1, z^2 \geq 0$$

The simplex method is performed with an additional rule :  $v'x + u'y = 0$  must hold throughout the procedure. This assures that at every point and for all  $i$ , no  $v_i$  and  $x_i$  can have nonzero value (cannot be in the basis) at the same time.

Wolfe's method exists in two forms: The short form that finds a minimum only if  $b = 0$  or  $A$  is positive definite and

the long form that can be applied without restrictions. The short form requires two phases in the procedure and the long form adds one more phase to the results obtained with the short form. In the first phase,  $r$  is set to 0, such that  $rb = 0$  (since in the short form, this guarantees a solution) and  $w$  is minimized (in fact, driven to zero), holding  $r$ ,  $u$  and  $v$  out of the basis.

The initial tableau is set:

$$w = q - Px - y$$

$$z^1 = 0 - 'x + v - P'u + z^2 - rb$$

and the simplex method is applied until  $\sum_{i=1}^n w_i = 0$ .

The second phase minimizes  $z^1$  and  $z^2$ . At this point, there is no need to differentiate between  $z^1$  and  $z^2$  since both must vanish.  $r$  must not enter the basis in this phase. Again, the simplex method is used until the sum of all entries in  $z^1$  and  $z^2$  is zero. This concludes the short form.

In the third phase, (in the long form)  $r$  is set to 1, and the objective is to obtain two successive tableaux in which the  $r$  values are slightly below and above 1. With each of these  $r$ 's there is associated a solution vector and a linear interpolation is used to obtain the solution vector corresponding to  $r=1$ , which is the problem solution.

## 4.2 THE METHOD OF DANTZIG

G.B. Dantzig [11] took the algorithm of Wolfe and developed a slightly more compact version of the method that works somewhat faster if  $A$  is either positive definite or semidefinite. The significant shortcoming is that it may not terminate at all if  $A$  is indefinite.

It uses the Kuhn-Tucker conditions (K<sup>T</sup>2) and the initial tableau is set as:

$$v = b + Ax + P'u$$

$$y = q - Px$$

An important distinction between two kinds of solutions must be stated: A tableau is said to be in standard form if  $v'x + u'y = 0$  holds. Otherwise it is in nonstandard form. The initial tableau is in standard form:  $v$  and  $y$  are basic with nonzero values and  $u$  and  $x$  are nonbasic, hence their values are zero. Whenever a tableau is in nonstandard form, the next iteration takes it back to the standard form.

The simplex method is performed with some special rules for interchanging the variables: If a tableau is in standard form, the variable to enter the basis is the  $x_k$  whose corresponding  $v_k$  (basic) has the largest positive value. If a tableau is in nonstandard form, the variable to enter the basis is the nonbasic  $v_j$  whose corresponding  $x_j$  is also nonbasic. The variable to be deleted from the basis is

selected with the same minimum ratio criterion used in the simplex method for linear programming [7,10].

The objective is to minimize  $\sum_{j=1}^n v_j$ . A standard tableau is optimum when the values of  $v$  in the basis are all negative: that is, no  $v$  applies for a new iteration. The  $x$  vector in this optimum tableau is the solution to the problem.

#### 4.3 EXPERIMENTATION

In Braitsch's study some preliminary experiments were carried out to help select the factors and levels for the primary experimentation. Listed below are the nine factors selected that may be important in determining how fast a quadratic programming algorithm will converge.

1. Algorithm : Refers to the algorithms of Wolfe, Dantzig, Beale and W.B. The following factors apply to each of the four algorithms.
2. Rule : Rule used in each algorithm for selecting the variable to enter the basis. In all cases but three, the variable entered was one that resulted in the largest rate of increase in the criterion function. The exceptions were WB, where no selection rule was necessary and the nonstandard form cases of Beale and Dantzig.

3. Size : Defined as the number of x-variables or the number of constraints in the P matrix, since squared P-matrices of size  $m \times m$  were used. Sizes of 15, 20 and 25 were used.
4. Negative b : The number of negative coefficients in the cost vector b. Levels of 3, 9 and 15 were used.
5. Rank : Number of linearly independent columns in the matrix A. Ranks of 3 and 15 were used.
6. Range : Range of the random numbers used to generate the P-matrix on the computer. The ranges (r) used were :  $-20 \leq r \leq 20$  and  $-180 \leq r < 180$
7. b-range : Range of the random numbers used to generate the b-vector. Ranges used were  $-20 \leq r \leq -1$  and  $-180 \leq r \leq -1$
8. q-range : Ranges for the random numbers in the q vector. Ranges used were  $1 \leq r \leq 20$  and  $1 \leq r \leq 180$
9. P-range : Only the single level of  $1 \leq r \leq 180$  was used for the random numbers that made up the P-matrix.

A complete factorial model was used where, for each of the four algorithms, all combination of factors (3) through (8) were run. The experiment was replicated ten times. The various factor combination were run by nesting a series of loops in the computer program, each loop corresponding to a factor. It was assumed that since the problem components

were randomly generated, further randomization would not be necessary. Each randomly generated problem was solved by all four algorithms.

#### 4.4 RESULTS AND ANALYSIS

TABLE 2  
MEAN NUMBER OF ITERATIONS FOR DIFFERENT LEVELS

FACTOR		DANTZIG	BFALF	W.B.	WOLFF
Size :	15	6.9	7.8	8.3	19.5
	20	7.5	8.7	8.0	21.6
	25	7.8	9.3	9.6	23.3
Negative p :	3	4.5	4.6	5.3	7.6
	9	7.8	8.8	9.4	20.9
	15	9.9	12.4	12.1	35.8
Rank :	3	7.3	8.6	8.9	21.6
	15	7.5	8.7	9.0	21.3
A-range :	[-20,20]	6.1	4.8	7.6	22.9
	[-180,180]	8.7	12.4	10.3	20.1
b-range :	[-20,-1]	7.9	10.1	9.5	20.6
	[-180,-1]	6.9	7.1	8.4	22.3
q-range :	[1,20]	7.2	7.7	8.6	22.5
	[1,180]	7.6	9.5	9.3	20.4
Algorithm means		7.4	8.6	8.9	21.5

Table 2 shows the average number of iterations taken at each level of all of the factors that were varied, for each of the algorithms.

The data in this table lead to the conclusions that the rank of the matrix has little effect, but the size factor is very significant. A number of the results for Dantzig's and Beale's algorithms are explainable in terms of the relative importance of the linear and nonlinear part of the objective function. It is clear that increasing the b-range or decreasing the A-range will emphasize the linear part of  $G(x)$ . Also, decreasing the q-range or increasing the P-range increases the importance of the linear part. This is true because these latter two actions tend to lower x-values which reduces the contribution of the quadratic part more than that of the linear part.

Average iterations decreased in the three cases mentioned above, that is increasing the b-range, decreasing the A-range and decreasing the q-range. This is reasonable, because in the extreme case when the quadratic portion is irrelevant and the problem becomes a linear programming problem, both methods reduce to the Simplex method which moves between adjacent extreme points in search of solution. In the nonlinear case, when the quadratic part is significant, the solution often lies in one of the faces in the feasible region rather than at an extreme point. This tends to add iterations because of the necessity of optimizing on faces. These extra iterations often include increasing the number of basic primal variables.

There is a major difference in the way Beale's and Dantzig's methods perform in the Linear Programming case. While Beale's method reduces to the Simplex method in the constraints portion, Dantzig's alternates between regular Simplex iterations and Dual Simplex iterations in two parts of the tableau, thus taking twice as many iterations as Beale's. As the quadratic part becomes more important, average iterations increase at a greater rate in Beale's algorithm than in Dantzig's. The major reason for this is that when we move from one face in the feasible region to another (by making a sign restricted variable nonbasic), it is necessary to replace all of Beale's free variables because these variables were used to optimize on the previous face.

Since each unit of size augmented adds a constraint and an x-variable, the new constraint adds another face to the feasible region which could raise iterations if the algorithm had to examine it on the way to the optimal solution.

Negative b is a very important factor in Dantzig's and Beale's methods. Each added negative b-coefficient adds a variable that can profitably enter in the first iteration for both algorithms. Although not all of the x-variables having negative b-coefficients will be in the optimal solution, the number will tend to increase with negative b

and iterations will be required to pivot them into the basis. This discussion does not apply to Wolfe's algorithm because it uses a different type of objective function than the other three algorithms. Algorithm WB generally requires fewer iterations because it uses only one artificial variable. The performance of this algorithm is difficult to explain, for although it is a modification of Wolfe's algorithm, it reacts to different problem types in the same way as Dantzig's.

#### 4.5 COMPUTER TIME AND STORAGE

It is difficult to get an accurate measure of the computer time required by the various algorithms because of the large difference that programming technique can cause. A more reliable measure might be based on the number of tableau elements that must be transformed at each iteration. In this way, we can take advantage of theoretical shortcuts for improving time efficiency that might normally be ignored by a programmer. These shortcuts include transforming only half of Beale's Objective tableau because it is symmetric, eliminating transformations in Dantzig's standard-form tableaux because of symmetry relations and reducing the

space that must be set aside due to free variables in Beale's algorithm by noting that the number of these free variables cannot exceed the rank of the A-matrix.

In the experimentation, approximately 70 percent of Dantzig's tableaux were in standard form. All tableaux transformations were assumed to take one unit of time, except for the transformations in Beale's objective tableau, which took about 1.6 units in computer tests. The time to transform variables added in Beale's simplex tableau due to free variables will be ignored because it is difficult to estimate, and judging from data available, probably would not affect the time estimate by more than a few percent at worst. Table 5 gives time and storage estimates for several combinations of problem parameters.

Timewise, Beale's and Dantzig's algorithms are superior, with Beale's algorithm tending to be preferable when the number of x-variables does not greatly exceed the number of constraints. This is due to the fact that Beale's tableaux are in terms of the nonbasic variables. In terms of storage, Beale's algorithm takes better advantage of situations where the P-matrix is nearly square or the rank of A is not a very large fraction of the number of rows or columns in A. The latter phenomenon is due to a reduction in space necessary

TABLE 3  
TIME AND STORAGE REQUIREMENTS

M	N	A-rank	Algorithm	Time	Storage
Z	Z	Z	BEALE	1.8	2.5
			DANTZIG	2.6	4
			W. B.	4	4
			WOLFF	4	4
Z	2Z	2Z	BEALE	5.2	8
			DANTZIG	6	9
			W. B.	9	9
			WOLFF	9	9
Z	4Z	4Z	BEALE	16.8	28
			DANTZIG	16.3	25
			W. B.	25	25
			WOLFF	25	25
Z	4Z	Z	BEALE	16.8	16
			DANTZIG	16.3	25
			W. B.	25	25
			WOLFF	25	25

With, N: number of variables, M: number of constraints  
Time: Time per iteration ( $Z^2$  units) and one unit =  $23 \times 10^{-6}$  secs. on the IBM 360-75  
Storage: Requirements of storage in memory ( $Z^2$  locations)

to hold the rows of variables driven out of the nonbasic set because of the introduction of free variables.

The results of this study suggest that in most of the cases, Beale's algorithm will generally take more iterations than Dantzig's, but the difference is not very big. But Beale's algorithm has some other advantages: It will find a local minimum even if A is not definite which may be of value in some applications; second, a redefinition of the

A-matrix, which when combined with a computer code using the product form of the inverse, can solve larger problems than Dantzig's algorithm, provided the rank of  $A$  is low.

## Chapter V

### THE INTERACTIVE PACKAGE

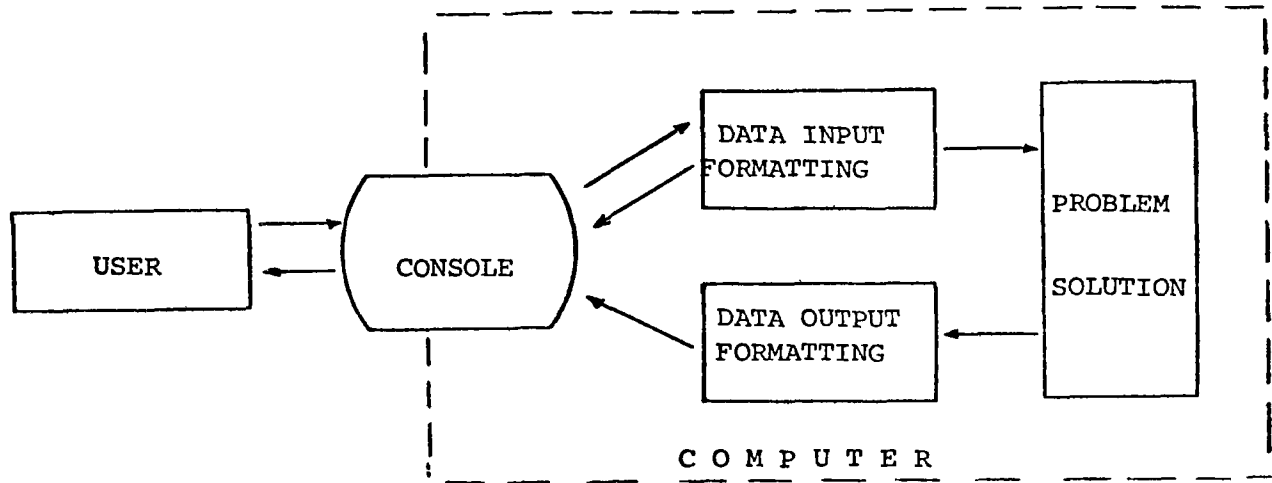
#### 5.1 STRUCTURE OF THE PACKAGE

The method of Beale was selected among the methods available after considering several important factors and, specifically, because

1. Beale's method does not require matrix  $A$  to be positive definite : it works also with  $A$  negative definite as well as in the semi-definite cases.
2. The tableaux used are small, giving economy in storage and computing time.
3. It uses the simplex method, thus is efficient.
4. The computational properties are acceptable or superior to the other methods compared.
5. Is the geometrically most illustrative among all methods for the QPP.
6. The procedure is relatively simple, thereby facilitating the programming task.

Figure 4 provides a general view of the package format. It consists of two major parts: A first part deals with the communication with the user, interchange of information and

Figure 4: PACKAGE INTERFACTION



the transformation of the data into a suitable format for both the user and the program itself. The second part is the procedure that actually solves the problem.

The execution of the package begins by giving the user a summary of its objective, the QPP is defined and some hints are given for putting the general objective function into matrix form. The user is informed about the procedure for communicating with the package.

No special format for the input data from the terminal is needed. The data entered by the user is checked and only the expected type of data is accepted. This is made with the aid of TAPE 99, a FORTRAN device that enables the reading of data with variable format.

The data received is displayed again on the terminal and the user is asked to check its accuracy. Corrections are permitted when errors are found. The tableau for the method of Beale is constructed from the data received and various parameters are initialized. The initial tableau is displayed if the user has chosen this output option.

The Beale algorithm is then initiated to solve the QPP, and again, according to the user's selected output option, intermediate tableaux and/or the final solution are displayed.

## 5.2 STORAGE ORGANIZATION

The procedure requires the matrix

$$\left[ \begin{array}{c|c} c & b/2 \\ \hline b/2 & \frac{1}{2}P \end{array} \right]$$

which is  $(n+1) \times (n+1)$  and symmetric (because  $P$  is symmetric). In order to save storage, only the upper triangular part of this matrix is stored. The  $m \times (n+1)$  matrix

$$\left[ q \mid -P \right]$$

is also created.

All 2-dimensional arrays required in the procedure are stored in a vector format which is created from the catenation of the rows of the matrix. For example, the 2x3 matrix P, where

$$P = \begin{bmatrix} P_{11} & P_{12} & P_{13} \\ P_{21} & P_{22} & P_{23} \end{bmatrix}$$

is stored as the 6-vector

$$v = \begin{bmatrix} P_{11} \\ P_{12} \\ P_{13} \\ P_{21} \\ P_{22} \\ P_{23} \end{bmatrix}$$

It then follows that the entry in row i, column j of P appears in v(k) where

$$k = (i-1)\alpha + j$$

and  $\alpha$  is the column dimension of P ( $\alpha = 3$  in this example). A slight modification of this process is used to store only the upper triangular part of symmetric matrices.

A means for distinguishing the variables that belong to the basis is required. For this purpose, the original nonbasic variables are always given the indices 1 through n

and the original basic variables the indices  $n+1$  through  $n+m$ . The free variables, referred to as "u" variables in chapter 3, are treated internally as x variables and they are assigned indices between  $n+m+1$  and  $2n+m$ . This was done to simplify their treatment and there is no risk of confusion for the user since these variables do not appear in the solution vector.

### 5.3 NOTATION AND DEFINITIONS

Throughout the programs, whenever possible, the same notation as that in chapter 3, was used.

- N The number of variables (user input); remains fixed during the execution of the procedure.
- M The number of constraints (user input); may vary during the procedure when free variables are added (each one yields a new constraint).
- VA A  $N*(N+1)/2$  - vector containing the upper triangular part of the A matrix in the quadratic objective function (user input).
- B Coefficient vector of the linear part of the objective function (user input).

C The constant value in the objective function (user input).

VP A M\*N - vector containing the entries of the constraint matrix P (user input).

Q M-vector with the right-hand side values of the constraints (user input).

EPS A small positive number corresponding to the width of the zero-band used in checks for zero at critical points during the solution (user input).

BASIC0 N-vector containing the indices of the nonbasic variables.

BASIC1 M-vector containing the indices of the basic variables.

LOS Vector containing the variables remaining in the objective function which are different from the free variables.

IAS Vector of the basic variables.

LIBRE Counts the free variables added to the problem.

ALIBRF Stores the number of free variables at a given point.

IPIV Stores the index of the pivot row.

KPIV Stores the index of the pivot column.

FL Set to n+1.

PICO Logical variable whose value is .TRUE. if the pivot is in the constraints part; otherwise, its value is .FALSE.

ISDONE Has a value 1 if the problem has been solved; otherwise it has a value of 0.

MAX Contains the maximum value in the first row of the Objective function tableau.

Y0 Contains the value of the minimum ratio  $a(i,0)/a(i,j)$  for which  $a(i,0)$  is nonnegative and  $a(i,j)$  is positive.

JC Set to  $N*(N+1)/2$ .

J1 Set to  $N*M$ .

SOL The vector which contains the problem solution.

DSPLAY The value assigned to this variable (by the user) serves to specify one of several available display options.

ITFP Counts the number of iterations that have been carried out.

#### 5.4 THE SUBPROGRAMS

The overall program is divided into twenty one FORTRAN subroutine subprograms. Each performs a specific task. These subroutines are linked and controlled by the Main program which serves as a driver routine.

#### 5.4.1 The Main program

This program calls subroutines DIALOG, INITIA, SETPTX, SHOWIT, and FUNIT. It interchanges information among them and transfers control from one to the other as the solution process is carried out.

#### 5.4.2 Subroutine DIALOG

Presents the package to the user, defines the problem and the variables, gives instructions for the use of the program and gets from the user the values of the parameters to be used in the problem, checks the validity of these values and if necessary accepts any correction suggested. The subroutines called from within DIALOG are READIT, READPM and MESSAG. This subprogram is called only by the Main program.

#### 5.4.3 Subroutine READIT

Reads one integer value from the terminal (specifically for N,M and DISPLAY). Analyzes the data received and sends messages to the user when necessary. This subprogram calls subroutine MESSAG and is called only by subroutine DIALOG.

#### 5.4.4 Subroutine READEM

This program reads a predetermined set of real values from the terminal. Transforms the data into double precision numbers and analyzes each input; accepts it or sends an error message when necessary. It is used to read the values of A, B, C, P, Q and FPS. This subprogram calls subroutine MFSSAG and is called only by subroutine DIALOG.

#### 5.4.5 Subroutine INITIA

It initializes the values of the parameters to be used in the program. Sets values for BASIC0, 'BASIC1, LOS, LAS, JC, J1, LIBF, ALIBF, LIBFE, RL and ITRF. This subprogram is called only by the Main program.

#### 5.4.6 Subroutine SETMTX

Transforms the data received from the user into a form suitable for use in Beale's method; i.e., it creates the matrix  $\bar{P}$  shown in section 3.1. It also creates the matrix  $[q | -P]$  for the constraints tableau. This subprogram is called only by the Main program.

#### 5.4.7 Subroutine SHOWIT

Displays information relating to the tableaux according to the value of DISPLAY. It is called at the end of each iteration to inform the user about the progress of the computations. If DISPLAY = 2, subroutines PROFUN and PRICCN are called. If DISPLAY = 1, the current solution is displayed. If DISPLAY = 0 control is returned to the calling routine. This subroutine is called by bc+1, the Main program and the subroutine RUNIT.

#### 5.4.8 Subroutine RUNIT

It is the subroutine that carries out Beale's algorithm. It calls subroutines MALIB, COLPIV, SETV, PATIC, CONACT, OFUACT, ADDONE, SHOWIT, CHANGE, DELETE and OPTIMO. It performs certain operations and transfers control to the subroutines, according to the results obtained. Controls the interaction of all the subroutines and identifies when the problem does not have a solution or when a solution has been found. This subprogram is called only by the Main program.

#### 5.4.9 Subroutine MALIB

This subroutine finds the largest value in the first row of the objective function tableau among the columns associated with the free variables. It is called each time a free variable is added to the problem. If MALIB finds any

positive value, this position defines the column of the pivot for the next iteration. It is called only by subroutine RUNIT.

#### 5.4.10 Subroutine COLPIV

It finds the column of the pivot for the next iteration. If MALIB has not found a positive value, COLPIV searches in the first row of the objective function tableau for the most negative number and defines that position to be the column of the pivot for the next iteration. If COLPIV doesn't find any negative value, we are at the optimum point. COLPIV is called only by subroutine PUNIT.

#### 5.4.11 Subroutine SETV

SETV sets the value of a variable V to be either 1 or -1, according to the sign of the element in the matrix  $\Gamma$  located at the intersection of its main diagonal and the pivot column. It is called only by subroutine PUNIT.

#### 5.4.12 Subroutine RATIO

This subroutine finds the Pivot row. To do this, it determines the least value of the ratio  $a(i,0)/a(i,j)$  where j is the column of the pivot, as i takes on its admissible values. If there are 2 different i's which yield the same minimum value, then this operation is repeated for all j's

different from the pivot column. This second step is necessary in order to avoid degeneracy in the next iteration. It is called only by subroutine RUNIT.

#### 5.4.13 Subroutine CONACT

It updates the tableau of the constraints. Performs the simplex procedure to interchange the variables determined by the pivot. This subroutine is called only by subroutine RUNIT.

#### 5.4.14 Subroutine OFUACT

It updates the tableau of the objective function. Performs the necessary computations to update the tableau at each iteration. It is called by subroutine RUNIT.

#### 5.4.15 Subroutine ADDONE

Adds one more free variable to the problem. This means also a new constraint to the system that goes to the simplex tableau. This subroutine is called only by subroutine RUNIT.

#### 5.4.16 Subroutine CHANGE

It performs the interchange of variables after each iteration and modifies the values of the parameters affected by this change. It is called by subroutine RUNIT.

#### 5.4.17 Subroutine DELETE

This subprogram deletes a constraint from the system, after the value of the free variable that originated it has reached a zero value in the objective function and a new constraint enters the system. It is called by the subroutine PUNIT.

#### 5.4.18 Subroutine PROFUN

It prints the tableau of the objective function. It reconstructs the matrix  $\Gamma$  and shows the result. Up to five columns are displayed at one time. If the tableau has more than five columns, they are successively displayed in groups of five until all have been shown. This subprogram is called by subroutines SHOWIT and OPTIMC.

#### 5.4.19 Subroutine PRICON

This subroutine prints the tableau of the constraints. It reconstructs the P matrix and as in PROFUN, only five columns are shown at a time. The remaining columns are displayed in subsequent tableaux. PRICON is called by subroutines SHOWIT and OPTIMC.

#### 5.4.20 Subroutine OPTIMO

It displays the problem solution. Prints a suitable message depending on whether ISDONE equals 0 or 1. If DISPLAY = 2, it calls PPOFUN and PRICON, otherwise, creates the solution vector SOI and displays it with the appropriate messages. It is called by subroutine FUNIT.

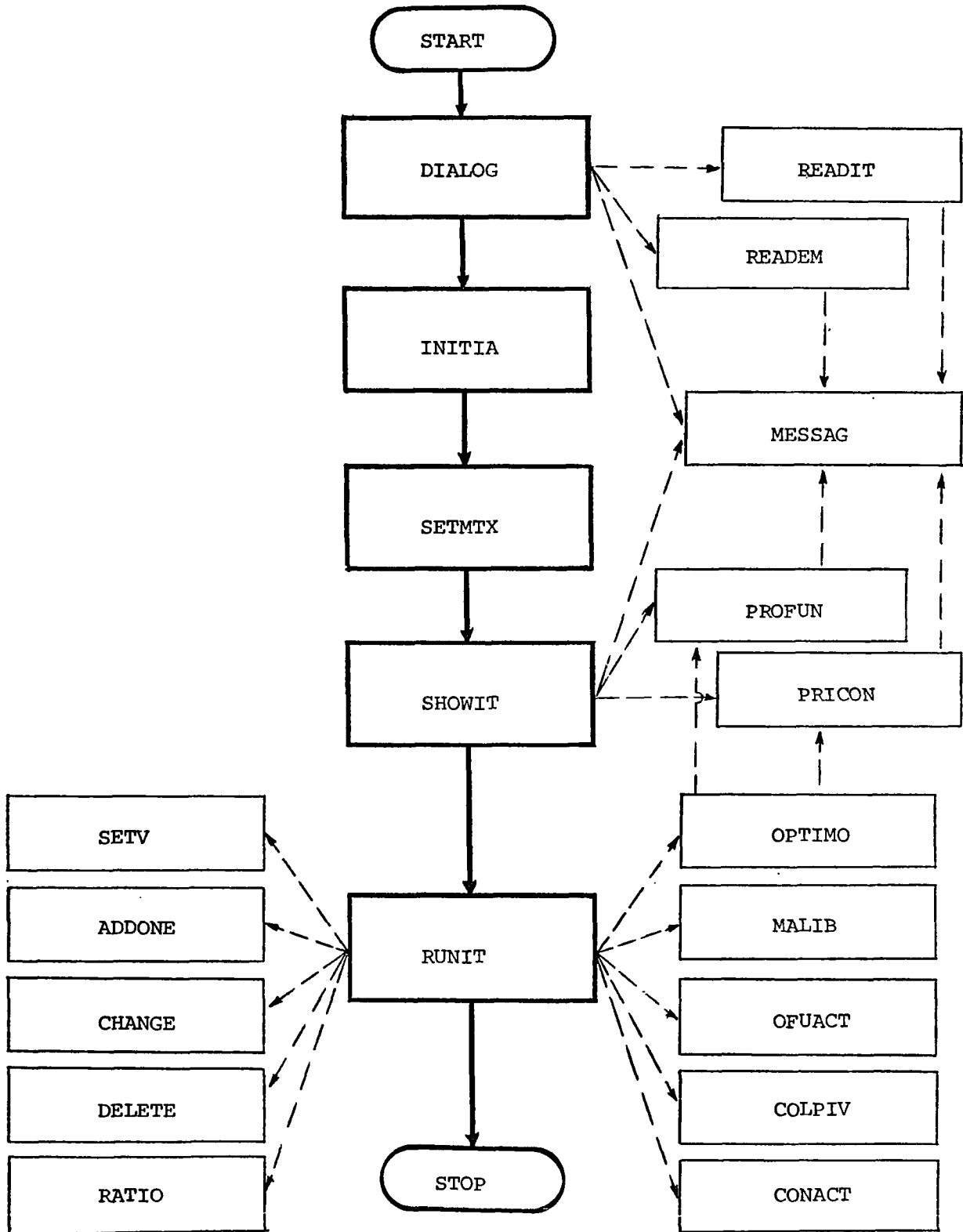
#### 5.4.21 Subroutine MESSAG

Prints messages for the user. It is called by subroutines DIALOG, PPOFUN, PRICON, PEADDM, PEADIT, SHOWIT and OPTIMO.

### 5.5 PROGRAM ASSEMBLY

Figure 5 shows how the routines are interconnected. The darker blocks represent the subprograms that are sequentially called within the main program. The dotted arrows represent subroutine calls.

Figure 5: THE PROGRAM AND THE SUBROUTINES



Chapter VI  
USING THE PACKAGE

6.1 CONNECTING TO THE PACKAGE

The package is to be used under VM-CMS. To run it, the user must have the following files in his disk

CUADRO EXEC  
SETB99 TXTLIB  
CUADRO FORTRAN

The file CUADRO EXEC contains:

```
GLOBAL TXTLIB FORTLIB
FORTRAN &1
FI 5 TERMINAL(RECFM F LRECL 80
FI 6 TERMINAL(RECFM F LRECL 80
CP SPOOL CONS START *
LOAD &1 (START
CP SPOOL CONS STOP CLOSE
&TYPE -----
&TYPE TYPE 'CUADRO PROBLEM' TO ENTER FILENAME FILETYPE
&TYPE -----
EXEC RECEIVE
&TYPE -----
```

```

&TYPE DO YOU WANT A HARD COPY OF THE INFORMATION DISPLAYED ?
&TYPE -----
&TYPE (<CR> IF NO, ELSE 'YES')
&READ VARS &ANSWER
&IF .&ANSWER NE . &GOTO -YES
&ERASE &1 PROBLEM
&EXIT
-YES PRINTOUT &1 PROBLEM

```

The file SETB99 TXTLIB (also known as "tape 99" or "buffer 99"), is a 132-byte array that exists in the region of core assigned to the programme. The space for this buffer is automatically created if the compiler detects the use of "tape unit 99".

The file CUADRO FORTRAN contains the source program of the package (see Appendix B).

To run the package, the user must type the command

```
cuadro_cuadro
```

## 6.2 PREPARING THE DATA

The package is organized to handle any problem that is in the following standard form

$$\text{minimize } \{ (1/2)x'Ax + b'x + c \mid Dx \leq q, x \geq 0 \}$$

To illustrate how a problem can be put into standard form, the example discussed in section 3.2 is taken as model. Suppose the problem is given as minimize G where

$$G = -6x_1 + 2x_1^2 - 2x_1x_2 + 2x_2^2$$

$$\text{Subject to } x_1 + x_2 \leq 2$$

$$x_1, x_2 \geq 0$$

It can be easily verified that this is equivalent to minimize G

where

$$G = (1/2) \begin{bmatrix} x_1 & x_2 \end{bmatrix} \begin{bmatrix} 4 & -2 \\ -2 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} -6 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} 0 \end{bmatrix}$$

$$\text{Subject to } \begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \leq \begin{bmatrix} 2 \end{bmatrix}, \quad x_1, x_2 \geq 0$$

The values assigned by the user (see Appendix A) are as follows

$$\begin{array}{rcc}
 N = 2 & & M = 1 \\
 VA = \begin{bmatrix} 4 \\ -2 \\ 4 \end{bmatrix} & b = \begin{bmatrix} -6 \\ 0 \end{bmatrix} & c = [0] \\
 VP = \begin{bmatrix} 1 \\ 1 \end{bmatrix} & & q = [2]
 \end{array}$$

Note that  $q$  must be strictly positive. If in a problem the standardized constraint part is not immediately available (by inspection), a basic starting solution must be searched. The technique discussed in section 2.3.1 can be used here, since the constraints part of the QPP has the same form as that of the linear programming problem.

### 6.3 MESSAGES

During execution of the package, various types of messages are transmitted to the user's terminal. These can be

1. Informative messages: where no answer is expected from the user; i.e. messages outlining the package, its outputs, etc.
2. Data requesting messages: where an answer (numerical data) is expected from the user; i.e. 'TYPE THE

VALUE OF N (I.E. THE NUMBER OF VARIABLES YOU HAVE)',  
'TYPE THE VALUE OF YOUR m-VECTOR Q. (THE RIGHT-HAND  
SIDE VALUE IN YOUR CONSTRAINTS)', etc.

3. Control messages: where an answer 'yes'(hit return)  
or 'no'(type no) is expected from the user, example  
'YOU HAVE n VARIABLES ? ', 'ARE THESE THE CORRECT  
VALUES OF YOUR m-VECTOR Q (THE RIGHT-HAND SIDE TERMS  
IN THE CONSTRAINTS) ? q q ... q ', etc.

4. Error messages: when the data is of a different kind  
than expected, the user is asked to input again,  
example

```
'-----> # ***** ERROR IN YOUR ENTRY : 12#456 ,  
INPUT AGAIN OR TYPE Q TO QUIT'  
'THE NUMBER OF CONSTRAINTS CANNOT BE NEGATIVE ! TRY  
AGAIN .. .  
etc.
```

#### 6.4 INTERPRETING THE RESULTS

If the value in the display option (DISPLAY) is 0, the package provides only the optimum value of the objective function together with the associated x-vector. If the choice is 1 the current value of the objective function at

each iteration is provided, from the starting point to the final one.

When the choice of display is 2, the two main tableaux are displayed at each iteration. An example of such a display is shown in table 4

Indeed, the form of the tableaux used in Peale's method contains complete information about the problem at any step. The element in the upper left corner the objective function tableau (i.e.  $g(1,1)$ ) contains the value of the objective function in that trial point. In the starting point, this corresponds to the value of  $c$ .

All the variables in the objective function tableau have a value of zero. All the variables in the constraint tableau have a nonzero value, given by the constant part in the constraint equation. However, if that variable has a subscript greater than  $N$ , then it is a slack variable (or a free variable) which has a zero coefficient in the objective function, hence it has no effect on the value of  $G$ .

The package contains a feature which permits the generation of a hard-copy of all information displayed on the screen during the computing session. This is achieved in the following way:

1. After the completion of the QPP solution, some information is displayed as the process of storing the solution file is progressing. The user is asked

TABLE 4  
AN EXAMPLE OF TABLE OUTPUT

---

ITERATION NUMBER 2

OBJECTIVE FUNCTION

	1	2	3	4	
	X 5	X 2	X 3	X 6	
	-0.20740D 02	0.24800D 01	-0.22400D 01	0.17900D 02	0.13400D 01
X 5	0.24800D 01	0.12000D 00	-0.20000D 00	0.40000D 00	0.0
X 2	-0.22400D 01	-0.20000D 00	0.53200D 01	-0.22000D 01	0.48000D 00
X 3	0.17900D 02	0.40000D 00	-0.22000D 01	0.45000D 01	-0.30000D 00
X 6	0.13400D 01	0.0	0.48000D 00	-0.30000D 00	0.22000D 00

CONSTRAINTS

	X 5	X 2	X 3	X 6	
X 1=	0.40000D 00	-0.20000D 00	0.0	-0.20000D 01	0.0
X 4=	0.60000D 00	0.0	-0.80000D 00	0.0	-0.20000D 00
X 7=	0.66667D 00	0.20000D 00	-0.20000D 00	0.10000D 01	0.20000D 00

THE CURRENT VALUE OF THE OBJECTIVE FUNCTION IS -20.74000000

THE VALUES OF THE VARIABLES ARE

X 1 = 0.40000000  
X 4 = 0.60000000  
X 7 = 0.66666700

ALL THE OTHERS ARE EQUAL TO ZERO

---

to type CUADFC PROGRAM as the FILENAME and FILETYPE

requested. If <CR> is made instead, no information is received.

2. The message 'DO YOU WANT A HARD-COPY OF THE INFORMATION DISPLAYED ?' (<CF> IF NOT, ELSE 'YES') is displayed and the desired answer is expected from the user.
3. If the answer is 'YES', the PRINTOUT command is sent, and the user is asked to type his batch account number and name.
4. If the answer is <CF> (no), the file CUADRO PROBLEMA is erased from the user's disk.

## Chapter VII

### CONCLUSIONS

The interactive computing environment is acknowledged as a far superior problem solving environment than the batch mode. The intent of the work outlined in this thesis has been to provide such an environment for solving the quadratic programming problem. The package which has been developed makes possible a convenient and straight forward solution to problems of this class and in particular simplifies the process of specifying the problem parameters. Very little technical expertise is required of the user. Note also that the package is able to handle the situation where  $M=0$ ; i.e. when the problem is unconstrained.

A highly modular organization has been followed in the development of the programme which should facilitate extensions or changes if these become required. The CUADRO package is a complete unit on its own. It could nevertheless be easily restructured to form part of some other package e.g. a general function minimization package. The calculations in the programme are carried out in double precision mode in order to minimize the influence of round-off errors. Inasmuch as FORTRAN IV was used as the programming language, extensive use of structured

programming concepts could not be employed; nevertheless an attempt was made to avoid excessive use of branching (GO TO) statements in order to improve the readability of the code.

At the present time, the dimensioning of various arrays in the program restricts the total number of constraints and variables to be less than 100. In order to handle larger problems, these arrays would have to be appropriately re-dimensioned.

## Appendix A

### AN EXAMPLE PROBLEM RUN WITH THE PACKAGE

```
cuadro cuadro
GLOBAL TXTLIB FORTLIB
FORTRAN CUADRO
FI 5 TERMINAL (RECFM F LRECL 80
FI 6 TERMINAL (RECFM F LRECL 80
CF SPOOL CONS STAFF *
LOAD CUADRO ( STAFF
EXECUTION BEGINS...
```

UNIVERSITY OF OTTAWA  
SYSTEMS SCIENCE  
1982

---

#### CUADRO

---

THIS PACKAGE USES THE METHOD OF BFGS TO SOLVE THE QUADRATIC PROGRAMMING PROBLEM, HAVING N VARIABLES AND M CONSTRAINTS ( $M \leq N$ ), GIVEN IN ITS STANDARD FORM :

SUBJECT TO : 
$$\begin{aligned} \text{MIN } G &= (1/2)X'AX + P'X + C \\ PX &\leq Q && \text{WITH } Q(J) > 0 \\ X(K) &\geq 0 && K = 1, 2, \dots, N \\ &&& J = 1, 2, \dots, M \end{aligned}$$

WHERE : C : THE CONSTANT VALUE OF YOUR OBJECTIVE FUNCTION  
B : THE VECTOR OF THE LINEAR PART IN THE O. P.  
A : AN NXN MATRIX CONTAINING THE QUADRATIC COEFFICIENTS OF THE O. P.  
P : AN NXM MATRIX WITH THE COEFFICIENTS OF THE CONSTRAINTS  
Q : AN M-VECTOR WITH THE RIGHT-HAND SIDE VALUES OF THE CONSTRAINTS

NOTE THAT MATRIX A IS A SYMMETRIC MATRIX WITH  $A(I, J) =$   
TWICE THE COEFFICIENT OF  $X(K)**2$  WHEN  $K=I=J$  (THE DIAGONAL OF THE MATRIX)  
COEF. OF  $X(I)X(J)$  WHEN I DIFFERENT FROM J (THE OFF-DIAGONAL ELEMENTS)

WHEN TYPING THE DATA REQUESTED, PLEASE RECALL :

- \* AT LEAST ONE BLANK (SPACE) MUST SEPARATE SUCCESSIVE NUMBERS
  - \* THE DECIMAL POINT IS NOT NECESSARY, UNLESS IT IS PART OF THE NUMBER
  - \* IF YOU WANT TO QUIT TYPE "Q" INSTEAD OF THE INFORMATION REQUESTED
  - \* WHEN ANSWERING TRUE OR FALSE,  
TYPE "NO" AND <CR> IF FALSE, OR  
<CR> (HIT RETURN) IF TRUE
-

TYPE THE VALUE OF N (I.E. THE NUMBER OF VARIABLES YOU HAVE )  
-3

YOUR VALUE OF N = -3 IS NOT POSITIVE !  
R\*START . . .

TYPE THE VALUE OF N (I.E. THE NUMBER OF VARIABLES YOU HAVE )  
2

TYPE THE VALUE OF M (I.E. THE NUMBER OF CONSTRAINTS YOU HAVE)  
3

THE STANDARDIZED PROBLEM REQUIRES  $M \leq N$  . . . .  
IN YOUR PROBLEM, A MAXIMUM OF 2 CONSTRAINTS (M) WILL BE ACCEPTED

TYPE THE VALUE OF M (I.E. THE NUMBER OF CONSTRAINTS YOU HAVE)  
1

THE QUADRATIC PART OF THE OBJECTIVE FUNCTION, BEING A SYMMETRIC MATRIX,  
ONLY THE UPPER-TRIANGULAR PART MUST BE TYPED

TYPE THE 2 VALUES OF THE ROW NUMBER 1  
4 -2

TYPE THE VERY LAST ELEMENT OF YOUR MATRIX  
5

TYPE THE 2-VECTOR B, (THE LINEAR PART OF YOUR OBJECTIVE FUNCTION)  
-6 0

TYPE THE VALUE OF C, I.E. THE CONSTANT PART IN YOUR OBJECTIVE FUNCTION  
0

-----> 0  
\*\*\*\*\* ERROR IN YOUR ENTRY :

0  
INPUT AGAIN OR TYPE Q TO QUIT  
0

TYPE THE 2 VALUES OF THE ROW NO. 1 IN YOUR CONSTRAINTS MATRIX P  
1 1.5

TYPE THE VALUE OF Q IN YOUR CONSTRAINT (I.E., THE RIGHT-HAND SIDE VALUE)  
2

THE VALUES IN THE ZERO-BAND ARE TAKEN AS ZERO IN CRITICAL POINTS  
OF THE PROBLEM.

TYPE THE VALUE OF THE ZERO-BAND FOR YOUR PROBLEM  
1.d-7

NOW, LET US CHECK THE CORRECTNESS OF THE DATA RECEIVED

TYPE "NO" AND <CR> IF FALSE, OR  
<CR> (HIT RETURN) IF TRUE

YOU HAVE 2 VARIABLES ?

YOU HAVE 1 CONSTRAINTS ?

ARE THESE THE CORRECT 2 ELEMENTS OF THE ROW NUMBER 1 OF MATRIX A  
IN YOUR OBJECTIVE FUNCTION ?

4.0000 -2.0000

IS 5.0000 THE LAST ELEMENT OF THE MATRIX IN YOUR OBJECTIVE FUNCTION ?  
10

TYPE THE VERY LAST ELEMENT OF YOUR MATRIX

4

IS 4.0000 THE LAST ELEMENT OF THE MATRIX IN YOUR OBJECTIVE FUNCTION ?

ARE THESE THE CORRECT VALUES OF YOUR 2-VECTOR B  
(THE LINEAR PART IN YOUR OBJECTIVE FUNCTION) ?

-6.000 0.0

THE CONSTANT VALUE C IN YOUR OBJECTIVE FUNCTION IS 0.0 ?

ARE THESE THE CORRECT 2 VALUES OF ROW NUMBER 1  
IN YOUR MATRIX P (CONSTRAINTS) ?

1.0000 1.5000

10

TYPE THE 2 VALUES OF THE ROW NO. 1 IN YOUR CONSTRAINTS MATRIX P

1 1

ARE THESE THE CORRECT 2 VALUES OF ROW NUMBER 1  
IN YOUR MATRIX P (CONSTRAINTS) ?

1.0000 1.0000

ARE THESE THE CORRECT VALUES OF YOUR 1-VECTOR Q  
(THE RIGHT-HAND SIDE TERMS IN YOUR CONSTRAINTS) ?

2.0000

THE WIDTH OF YOUR ZERO-BAND IS 0.0000001000 ?

SELECT ONE OF THE FOLLOWING OPTIONS :

0 >ONLY THE FINAL ANSWER WILL BE SHOWN

1 >THE VALUES OF THE OBJECTIVE FUNCTION AND THE BASIC VARIABLES WILL BE SHOWN  
AT EACH ITERATION

2 >THE COMPLETE TABLEAU FOR THE METHOD OF BRUTE WILL BE DISPLAYED AT EACH  
ITERATION

TYPE THE NUMBER OF THE OPTION YOU CHOOSE

2

ORIGINAL TABLEAU

OBJECTIVE FUNCTION

	1	2
	Y 1	X 2
0.0	-0.30000D 01	0.0
1 -0.30000D 01	0.20000D 01	-0.10000D 01
2 0.0	-0.10000D 01	0.20000D 01

CONSTRAINTS

	X 1	X 2
3= 0.20000D 01	-0.10000D 01	-0.10000D 01

THE CURRENT VALUE OF THE OBJECTIVE FUNCTION IS 0.0

THE VALUES OF THE VARIABLES ARE

X 3 = 2.00000000

ALL THE OTHERS ARE EQUAL TO ZERO

---

ITERATION NUMBER 1

OBJECTIVE FUNCTION

	1	2
	X 4	X 2
-0.45000D 01	0.0	-0.15000D 01
X 4 0.0	0.50000D 00	0.0
X 2 -0.15000D 01	0.0	0.15000D 01

CONSTRAINTS

	X 4	X 2
X 3= 0.50000D 00	-0.50000D 00	-0.15000D 01
X 1= 0.15000D 01	0.50000D 00	0.50000D 00

THE CURRENT VALUE OF THE OBJECTIVE FUNCTION IS -4.50000000

THE VALUES OF THE VARIABLES ARE

X 3 = 0.50000000  
X 1 = 1.50000000

ALL THE OTHERS ARE EQUAL TO ZERO

---

ITERATION NUMBER 2

OBJECTIVE FUNCTION

1 2

X 4 X 3

	-0.53333D 01	0.33333D 00	0.66667D 00
X 4	0.33333D 00	0.66667D 00	0.33333D 00
X 3	0.66667D 00	0.33333D 00	0.66667D 00

CONSTRAINTS

X 4 X 3

X 2=	0.33333D 00	-0.33333D 00	-0.66667D 00
X 1=	0.16667D 01	0.33333D 00	-0.33333D 00

THE CURRENT VALUE OF THE OBJECTIVE FUNCTION IS -5.33333333

THE VALUES OF THE VARIABLES ARE

X 2 =	0.33333333
X 1 =	1.66666667

ALL THE OTHERS ARE EQUAL TO ZERO

---

ITERATION NUMBER 3

OBJECTIVE FUNCTION

	1	2
	X 4	X 3
	-0.55000D 01	0.0
X 4	0.0	0.15000D 01
X 3	0.50000D 00	0.0

CONSTRAINTS

	X 4	X 3
X 2 =	0.50000D 00	-0.50000D 00
X 1 =	0.15000D 01	0.50000D 00

THE CURRENT VALUE OF THE OBJECTIVE FUNCTION IS -5.50000000

THE VALUES OF THE VARIABLES ARE

X 2 =	0.50000000
X 1 =	1.50000000

ALL THE OTHERS ARE EQUAL TO ZERO

---

SOLUTION

OBJECTIVE FUNCTION

	1	2
	X 4	X 3
-0.55000D 01	0.0	0.50000D 00
X 4 0.0	0.15000D 01	0.0
X 3 0.50000D 00	0.0	0.50000D 00

CONSTRAINTS

	X 4	X 3
X 2= 0.50000D 00	-0.50000D 00	-0.50000D 00
X 1= 0.15000D 01	0.50000D 00	-0.50000D 00

THE OPTIMAL IS -5.50000000

THE VALUES OF THE VARIABLES ARE

X 2= 0.50000000  
X 1= 1.50000000

ALL THE OTHERS ARE EQUAL TO ZERO

OSCAR MANLIOU, 1982

CP SPOOL CONS STOP CLOSE  
CON FILE 0463 TO VWPSI

-----  
TYPE 'CUADRO PROBLEM' TO ENTER FILENAME FILETYPE  
-----

EXFC RECEIVE

PDR file 463 unnamed

Enter FILENAME FILETYPE for READ. <cr>=no RECEIVE

cuadro problem

RECORD LENGTH IS '132' BYTES.

File CUADRO PROBLEM A RECEIVED from VWPSI

-----  
DO YOU WANT A HARD COPY OF THE INFORMATION DISPLAYED ?  
-----

(<cr> IF NOT, ELSE YES)

yes

PRINTOUT CUADRO PROBLEM

batch account number vwpslxxx

programmer name oscar

PUN FILE 0464 TO OSVS1 COPY=0001 PFCDS=000318 NO HOLD

R; T=3.38/4.45 13:06:16

## Appendix B

### THE FORTRAN PROGRAMS

CALL DIALOG	CUB00010
CALL INITIA	CUB00020
CALL SFTMTX	CUB00030
CALL SHOWIT	CUB00040
CALL PUNIT	CUB00050
STOP	CUB00060
END	CUB00070
	CUB00080
	CUB00090
SUBROUTINE PUNIT	CUB00100
IMPLICIT REAL*8 (A-H,O-Z)	CUB00110
INTEGER BASICO, BASIC1, FL, ALTBFF, F, T, V	CUB00120
REAL*8 MAX	CUB00130
LOGICAL PICO	CUB00140
COMMON /NUMERC/ N, M	CUB00150
COMMON /SUBJECT/ P (200)	CUB00160
COMMON /LIBFFS/ LIBF (50), LIPFE, ALIBFF (50)	CUB00170
COMMON /BASICA/ BASICO (50), BASIC1 (50)	CUB00180
COMMON /VALUES/ EPS, MIC, TTF, DISPLAY, KIM	CUB00190
COMMON /MAXIMO/ MAX	CUB00200
COMMON /MINIMO/ YC	CUB00210
COMMON /PARAME/ V, T, F	CUB00220
COMMON /PIVOTE/ IPIV, KPIV	CUB00230
COMMON /KNUMBER/ KCO, KVA	CUB00240
COMMON /LENGTH/ FL, JC, J1	CUB00250
COMMON /LOGICO/ PICO	CUB00260
COMMON /FINAL / ISDONE	CUB00270
ISDONE = 1	CUB00280
1110 IF (LIBRE.EQ.0) GO TO 1120	CUB00290
CALL MALIB	CUB00300
IF (DABS (MAX).GT.EPS) GO TO 1200	CUB00310
1120 IF (KVA.LE.0) GO TO 1130	CUB00320
CALL COLPIV	CUB00330
IF (YO.LT.-EPS) GO TO 1200	CUB00340
1130 ISDONE=1	CUB00350
GO TO 1990	CUB00360
1200 CALL SFTV	CUB00370
CALL PATIO	CUB00380
IF (IPIV.NE.0.OF.MAX.NE.0) GO TO 1300	CUB00390
ISDONE=0	CUB00400
GO TO 1990	CUB00410
1300 IF (IPIV.EQ.0.OF.YC.GT.DABS (MAX)) GO TO 1400	CUB00420
I=BASICO (KPIV)	CUB00430
BASICO (KPIV)=BASIC1 (IPIV)	CUB00440
BASIC1 (IPIV)=I	CUB00450
IE=R	CUB00460

P=1	CUA00470
CALL CONACT	CUA00480
IF (ISDONE.EQ.0) GO TO 1990	CUA00490
P=I <sup>F</sup>	CUA00500
PICO=.TRUE.	CUA00510
CALL OFUACT	CUA00520
GO TO 1600	CUA00530
1400 ALIBFE (KPIV) =ALIBFF (KPIV) +1	CUA00540
T=0	CUA00550
F=0	CUA00560
1410 IF (LIBFE.EQ.0) GO TO 1500	CUA00570
DO 16 I=1,LIBFE	CUA00580
IF (LIBF (I) .EQ. KPIV) GO TO 1430	CUA00590
16 CONTINUE	CUA00600
GO TO 1500	CUA00610
1430 CALL ADDONE	CUA00620
CALL CONACT	CUA00630
IF (ISDONE.EQ.0) GO TO 1990	CUA00640
PICO=.FALSE.	CUA00650
CALL OFUACT	CUA00660
CALL SHOWIT	CUA00670
GO TO 1110	CUA00680
1500 CALL CHANGE	CUA00690
GO TO 1410	CUA00700
1600 IF (BASIC1 (IPIV) .LE. N+MIO) GO TO 1650	CUA00710
CALL DELFTE	CUA00720
1650 CALL SHOWIT	CUA00730
GO TO 1110	CUA00740
1990 CALL OPTIMO	CUA00750
RETURN	CUA00760
END	CUA00770
	CUA00780
	CUA00790
	CUA00800
SUBROUTINE INITIA	CUA00810
IMPLICIT REAL*8 (A-H,O-Z)	CUA00820
INTEGER BASICO,BASIC1,ALIBFE,FL	CUA00830
COMMON /NUMERO/ N,M	CUA00840
COMMON /BASICA/ BASICO (50) ,BASIC1 (50)	CUA00850
COMMON /LIBFES/ LIBR (50) ,LIBFF,ALIBFF (50)	CUA00860
COMMON /KNUMBER/ KCO,KVA	CUA00870
COMMON /LENGTH/ FL,JC,J1	CUA00880
COMMON /VALUES/ EPS,MIO,ITEP,DSPLAY,KTM	CUA00890
COMMON /OTROS / LCS (50) ,LAS (50)	CUA00900
JC= ( (N+1) * (N+2) ) /2	CUA00910
J1= (M+N+2) * (N+1)	CUA00920
EPS = DABS (EPS/2.D0)	CUA00930
NI=N+1	CUA00940
DO 10 I=1,NI	CUA00950
BASICO (I) =I	CUA00960
10 CONTINUE	CUA00970
DO 13 K=1,N	CUA00980
LOS (K) =K	

```

        ALIBFF (K) = 0
        LIBP (K) = K
13 CONTINUE
IF (M.EQ.0) GO TO 1102
DO 14 I=1,M
    LAS (I) = I
    BASIC1 (I) = N+I
14 CONTINUE
1102 KVA=N
    KCO=M
    IL=N+1
    KIM=N+M
    ITEP=0
    LIBPF=0
    MIO=M
    RETURN
END

```

```

SUBROUTINE SETMTX
IMPLICIT REAL*8 (A-H,O-Z)
DIMENSION G(200),P(100)
COMMON /NUMERO/ N,M
COMMON /SUBJECT/ P(200),Q(50)
COMMON /FUNCN/ A(100),B(50),C
KM=(N+1)*(N+2)/2
LM=(N+1)*N/2
H(1) = C
JH = 2
DO 15 I=1,N
    H(JH) = B(I)/2
    JH = JH + 1
15 CONTINUE
DO 16 J=1,LM
    H(JH)=A(J)/2
    JH=JH+1
16 CONTINUE
DO 17 K=1,KM
    A(K)=H(K)
17 CONTINUE
IT=(M+1)*(N+1)
NI=N+1
DO 11 K=1,NI
    G(K)=0.D0
11 CONTINUE
IN=N
I1=1
DO 19 J=1,M
    NI=NI+1
    G(NI)=Q(J)
    DO 20 IJ=I1,IN

```

```

CUA00990
CUA01000
CUA01010
CUA01020
CUA01030
CUA01040
CUA01050
CUA01060
CUA01070
CUA01080
CUA01090
CUA01100
CUA01110
CUA01120
CUA01130
CUA01140
CUA01150
CUA01160
CUA01170
CUA01180
CUA01190
CUA01200
CUA01210
CUA01220
CUA01230
CUA01240
CUA01250
CUA01260
CUA01270
CUA01280
CUA01290
CUA01300
CUA01310
CUA01320
CUA01330
CUA01340
CUA01350
CUA01360
CUA01370
CUA01380
CUA01390
CUA01400
CUA01410
CUA01420
CUA01430
CUA01440
CUA01450
CUA01460
CUA01470
CUA01480
CUA01490
CUA01500

```

	NI=NI+1	CUA01510
	G(NI)=-P(IJ)	CUA01520
20	CONTINUE	CUA01530
	I1=IN+1	CUA01540
	IN=IN+N	CUA01550
19	CONTINUE	CUA01560
	DO 21 K=1,IT	CUA01570
	P(K)=G(K)	CUA01580
21	CONTINUE	CUA01590
	RETURN	CUA01600
	END	CUA01610
		CUA01620
		CUA01630
	SUBROUTINE SETV	CUA01640
	IMPLICIT REAL*8 (A-H,O-Z)	CUA01650
	REAL*8 MAX	CUA01660
	INTEGER V	CUA01670
	COMMON /PIVOTE/ IPIV,KPIV	CUA01680
	COMMON /VALUES/ EPS	CUA01690
	COMMON /NUMERO/ N,M	CUA01700
	COMMON /OFUNCN/ A(100)	CUA01710
	COMMON /MAXIMO/ MAX	CUA01720
	COMMON /PAFAME/ V	CUA01730
	KI=KPIV*(N+1)-((KPIV-1)*KPIV)/2+1	CUA01740
	MAX=0.D0	CUA01750
	IF(A(KI).GT.EPS) MAX=A(KPIV+1)/A(KI)	CUA01760
	V=-1	CUA01770
	IF(MAX.GT.0.D0) V=1	CUA01780
	RETURN	CUA01790
	END	CUA01800
		CUA01810
		CUA01820
	SUBROUTINE ADIONE	CUA01830
	IMPLICIT REAL*8 (A-H,O-Z)	CUA01840
	INTEGER T,RL,Z	CUA01850
	COMMON /PIVOTE/ IPIV,KPIV	CUA01860
	COMMON /LENGTH/ FL	CUA01870
	COMMON /NUMERO/ N,M	CUA01880
	COMMON /PAFAME/ V,T	CUA01890
	COMMON /SUBJCT/ P(200)	CUA01900
	COMMON /OFUNCN/ A(100)	CUA01910
	Z=0	CUA01920
	KAP=KPIV+1	CUA01930
	DO 17 K=1,KAP	CUA01940
	JO=K-1	CUA01950
	KU=T*FL+JO+1	CUA01960
	AJO=KPIV+Z+1	CUA01970
	P(KU)=A(AJO)	CUA01980
17	Z=Z+P-JO	CUA01990
	NIL=N-KPIV	CUA02000
	IF(NIL.LT.1)GO TO 1450	CUA02010
	DO 18 K=1,NIL	CUA02020

```

      KC=T*FL+(KPIV+K)+1
      KOJ=KPIV*(N+1)-(KPIV*(KPIV-1))/2+K+1
18 P(KO)=A(KOJ)
1450 IPIV=M
      RETURN
      END

```

```

SUBROUTINE CHANGE
IMPLICIT REAL*8 (A-H,O-Z)
INTEGER T,P,BASIC1,BASIC0
COMMON /KNUMTE/ KCO,KVA
COMMON /PARAME/ V,T,F
COMMON /LIBRES/ LIBR(50),LIBRT
COMMON /PIVOTE/ IPTV,KPIV
COMMON /NUMEPC/ N,M
COMMON /BASICA/ BASIC0(50),BASIC1(50)
COMMON /OTPOS / LOS(50),LAS(50)
COMMON /VALUES/ EPC,MIC
LIBRE=LIBRT+1
LIBR(LIBRE)=KPIV
BASIC1(M+1)=BASIC0(KPTV)
BASIC0(KPIV)=N+MIO+KPIV
LAS(M+1)=M+1
KCO=M+1
IF(KVA.LT.1)GO TO 1540
DO 19 K=1,KV
  IF(LOS(K).EQ.KPIV)T=K
19 CONTINUE
  KVA=KVA-1
1540 IF(KVA.LT.1.OPT.LE.0)GO TO 1560
  DO 20 K=T,KVA
    LOS(K)=LOS(K+1)
  20 CONTINUE
1560 M=M+1
      T=M
      F=1
      RETURN
      END

```

```

SUBROUTINE DELETE
IMPLICIT REAL*8 (A-H,O-Z)
INTEGER FL,ALIBRT,BASIC1
COMMON /NUMEPC/ N,M
COMMON /KNUMEI/ KCC,KVA
COMMON /BASICA/ KBA0(50),BASIC1(50)
COMMON /PIVOTE/ IPIV,KPTV
COMMON /LIBRES/ LIBR(50),LIBRT,ALIBRT(50)
COMMON /OTPOS / LOS(50)
COMMON /LENGTH/ FL
COMMON /SUBJCT/ P(200)

```

```

CUA02030
CUA02040
CUA02050
CUA02060
CUA02070
CUA02080
CUA02090
CUA02100
CUA02110
CUA02120
CUA02130
CUA02140
CUA02150
CUA02160
CUA02170
CUA02180
CUA02190
CUA02200
CUA02210
CUA02220
CUA02230
CUA02240
CUA02250
CUA02260
CUA02270
CUA02280
CUA02290
CUA02300
CUA02310
CUA02320
CUA02330
CUA02340
CUA02350
CUA02360
CUA02370
CUA02380
CUA02390
CUA02400
CUA02410
CUA02420
CUA02430
CUA02440
CUA02450
CUA02460
CUA02470
CUA02480
CUA02490
CUA02500
CUA02510
CUA02520
CUA02530
CUA02540

```

BASIC1 (IPIV)=BASIC1 (M)	CUA02550
KCO=KCO-1	CUA02560
KVA=KVA+1	CUA02570
LOS (KVA)=KPIV	CUA02580
LIBFE=LIBRF-1	CUA02590
IF (LIBFE.LT.1) GO TO 1640	CUA02600
DO 21 K=1,LIBFE	CUA02610
IF (LIBR (K) .EQ. KPIV) GO TO 1630	CUA02620
21 CONTINUE	CUA02630
GO TO 1640	CUA02640
1630 DO 22 I=K,LIBFE	CUA02650
LIBF (I)=LIBR (I+1)	CUA02660
22 CONTINUE	CUA02670
1640 NI=N+1	CUA02680
KIE=IPIV*PL+1	CUA02690
KIA=M*PL+1	CUA02700
DO 23 K=1,NI	CUA02710
KII=K-1	CUA02720
KIU=KIP+KII	CUA02730
KIO=KIA+KII	CUA02740
P (KIU) =P (KIO)	CUA02750
23 CONTINUE	CUA02760
M=M-1	CUA02770
ALIBR (KPIV) =ALIBR (KPIV) -1	CUA02780
RETURN	CUA02790
END	CUA02800
	CUA02810
	CUA02820
SUBROUTINE PROFUN	CUA02830
IMPLICIT REAL*8 (A-H,O-Z)	CUA02840
DIMENSION X (100)	CUA02850
INTEGER B0	CUA02860
COMMON /OFUNCN/ A (100)	CUA02870
COMMON /NUMERO/ N	CUA02880
COMMON /BASICA/ B0 (50)	CUA02890
COMMON /VALUES/ EPS,MIC,ITFF	CUA02900
IF (ITFF.EQ.0) CALL MESSAG (15)	CUA02910
L=2*N+1	CUA02920
NI=N+1	CUA02930
NN=N+1	CUA02940
IF (NI.GT.5) NI=5	CUA02950
NO=NI-1	CUA02960
WRITE (6,214) (K,K=1,NO)	CUA02970
WRITE (6,215) (B0 (K),K=1,NO)	CUA02980
DO 216 I=1,NN	CUA02990
DO 217 K=1,NI	CUA03000
IF (K.GT.I) GO TO 219	CUA03010
K1= (L-K+1) * (K-1) /2+I	CUA03020
GO TO 217	CUA03030
219        K1= (L-I+1) * (I-1) /2+K	CUA03040
217        X (K) =A (K1)	CUA03050
IF (I.NE.1) GO TO 280	CUA03060

240	WRITE(6,220) (X(K),K=1,NI)	CDA03070
	GO TO 216	CDA03080
280	WRITE(6,222) B0(I-1), (X(K),K=1,NI)	CDA03090
216	CONTINUE	CDA03100
285	IF(NO.EQ.N) RETURN	CDA03110
	NI=NO+1	CDA03120
	NO=N	CDA03130
	WRITE(6,214) (K,K=NI,NO)	CDA03140
	WRITE(6,215) (B0(K),K=NI,NO)	CDA03150
	DO 226 I=1,NN	CDA03160
	DO 227 K=NI,NO	CDA03170
	K1=K+1	CDA03180
	IF(K1.LT.I) GO TO 228	CDA03190
	KI=(I-I+1)*(I-1)/2+K1	CDA03200
	GO TO 227	CDA03210
228	KI=(I-K1+1)*(K1-1)/2+I	CDA03220
227	X(K)=A(KI)	CDA03230
	IF(I.EQ.1) WRITE(6,230) (X(J),J=NI,NO)	CDA03240
	IF(I.EQ.1) GO TO 226	CDA03250
	WRITE(6,232) B0(I-1), (X(J),J=NI,NO)	CDA03260
226	CONTINUE	CDA03270
250	RETURN	CDA03280
214	FORMAT(/,5X,'OBJECTIVE FUNCTION',/,11X,4(12X,I2))	CDA03290
215	FORMAT(/,23X,4('X',I2,11X))	CDA03300
220	FORMAT(/,3X,5(2X,D12.5))	CDA03310
222	FORMAT(/,1X,'X',I2,1X,D12.5,4(2X,D12.5))	CDA03320
230	FORMAT(/,17X,4(2X,D12.5))	CDA03330
232	FORMAT(/,1X,'X',I2,13X,4(2X,D12.5))	CDA03340
	END	CDA03350
		CDA03360
		CDA03370
	SUBROUTINE PPICON	CDA03380
	IMPLICIT REAL*8 (A-H,O-Z)	CDA03390
	INTEGER B0,B1,PL	CDA03400
	COMMON /NUMERO/ N,M	CDA03410
	COMMON /BASIC/ B0(50),B1(50)	CDA03420
	COMMON /SUBJECT/ P(200)	CDA03430
	COMMON /LENGTH/ PL	CDA03440
	N1=N+1	CDA03450
	M1=M+1	CDA03460
	IF(N1.GT.5) N1=5	CDA03470
	NO=N1-1	CDA03480
	CALL MESSAG(16)	CDA03490
	WRITE(6,307) (B0(K),K=1,NO)	CDA03500
	DO 306 I=2,M1	CDA03510
	KI=(I-1)*PL+1	CDA03520
	KC=KI+NO	CDA03530
	WRITE(6,309) B1(I-1), (P(K),K=KI,KC)	CDA03540
306	CONTINUE	CDA03550
	IF(NO.EQ.N) RETURN	CDA03560
320	NI=NO+1	CDA03570
	NO=N	CDA03580

IF (NO-NI.GT.5) NO=NI+5	CUA03590
WRITE(6,307) (B0(K),K=NI,NO)	CUA03600
DO 310 I=2,M1	CUA03610
K1=(I-1)*RI+NI+1	CUA03620
K2=K1+NC-NI	CUA03630
WRITE(6,311) B1(I-1), (P(K),K=K1,K2)	CUA03640
310 CONTINUE	CUA03650
IF (NO.EQ.N) RETURN	CUA03660
GO TO 320	CUA03670
307 FORMAT(/,23X,4('X',I2,11Y))	CUA03680
309 FORMAT(/,1X,'X',I2,'=',D12.5,4(2X,D12.5))	CUA03690
311 FORMAT(/,1X,'X',I2,13X,4(2X,D12.5))	CUA03700
END	CUA03710
	CUA03720
	CUA03730
	CUA03740
SUBROUTINE MALIB	CUA03750
IMPLICIT REAL*8 (A-H,O-Z)	CUA03760
REAL*8 MAX	CUA03770
COMMON /OFUNCN/ A(100)	CUA03780
COMMON /LIBRES/ LIBF(50),LIBRF	CUA03790
COMMON /NUMEFC/ N,M	CUA03800
COMMON /PIVOTE/ KDUM,KPIV	CUA03810
COMMON /MAXIMC/ MAX	CUA03820
KPIV=LIBF(1)	CUA03830
KOH=LIBF(1)+1	CUA03840
MAX=A(KOH)	CUA03850
IF (LIBFE.LT.2) RETURN	CUA03860
DO 402 K=2,LIBFE	CUA03870
IF (MAX.EQ.0) GO TO 402	CUA03880
KOH=LIBF(K)+1	CUA03890
IF (A(KOH).EQ.0) GO TO 402	CUA03900
IF (DABS(MAX).GT.DABS(A(KOH))) GO TO 402	CUA03910
KPIV=LIBF(K)	CUA03920
MAX=A(KOH)	CUA03930
402 CONTINUE	CUA03940
RETURN	CUA03950
END	CUA03960
	CUA03970
	CUA03980
	CUA03990
SUBROUTINE OFUACT	CUA04000
IMPLICIT REAL*8 (A-H,O-Z)	CUA04010
INTEGER FL,I,Z,Z1	CUA04020
LOGICAL PICC	CUA04030
COMMON /SUBJCT/ P(200)	CUA04040
COMMON /OFUNCN/ A(100)	CUA04050
COMMON /LENGTH/ RI	CUA04060
COMMON /PIVOTE/ IP,KP	CUA04070
COMMON /NUMEFC/ N,M	CUA04080
COMMON /LOGICO/ PICO	CUA04090
IN=N+1	CUA04100
DO 502 IP=1,IN	
J=IP-1	

	IF (J.EQ.KP) GO TO 502	CUA04110
	Z=0	CUA04120
	IF (J.GT.KP) Z1=KP-1	CUA04130
	IF (J.LT.KP) Z1=J	CUA04140
	IZ1=Z1+1	CUA04150
	DO 504 IS=1, IZ1	CUA04160
	K=IS-1	CUA04170
	KH=J+Z+1	CUA04180
	KH1=IP*FL+J+1	CUA04190
	KH0=KP+Z+1	CUA04200
	A (KH) =A (KH) +P (KH1) *A (KH0)	CUA04210
504	Z=Z+N-K	CUA04220
	T=Z+KP	CUA04230
502	CONTINUE	CUA04240
506	IKP=KP+1	CUA04250
	IF (IKP.GT.N) GO TO 530	CUA04260
	DO 510 K=IKP, N	CUA04270
	DO 512 J=K, N	CUA04280
	KH=Z+N-KP+J+1	CUA04290
	KH1=IP*RL+J+1	CUA04300
	KH0=T+K-KP+1	CUA04310
512	A (KH) =A (KH) +P (KH1) *A (KH0)	CUA04320
510	Z=Z+N-K	CUA04330
	DO 522 J=IKP, N	CUA04340
	KH=T+J-KP+1	CUA04350
	KH1=IP*PL+J+1	CUA04360
522	A (KH) =A (KH) +P (KH1) *A (T+1)	CUA04370
530	Z1=0	CUA04380
	DO 532 IR=1, KP	CUA04390
	J=IR-1	CUA04400
	Z=0	CUA04410
	KH=KP+Z1+1	CUA04420
	KH1=IP*PL+J+1	CUA04430
	SUMA=A (KH) +P (KH1) *A (T+1)	CUA04440
	DO 534 IS=1, IF	CUA04450
	K=IS-1	CUA04460
	KH=IP*FL+K+1	CUA04470
	KH1=IB+Z	CUA04480
	A (KH1) =A (KH1) +P (KH) *SUMA	CUA04490
534	Z=Z+N-K	CUA04500
532	Z1=Z1+N-J	CUA04510
	IF (IKP.GT.N) GO TO 550	CUA04520
	DO 542 J=IKP, N	CUA04530
	Z=0	CUA04540
	IP=J+1	CUA04550
	DO 542 IS=1, IR	CUA04560
	K=IS-1	CUA04570
	IF (K.NE.KP) GO TO 546	CUA04580
	Z=Z+1	CUA04590
	GO TO 542	CUA04600
546	KH=J+Z	CUA04610
	KH0=T+J-KP+1	CUA04620

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          KH1=IP*FL+K+1
          A (KH) =A (KH) +P (KH1) *A (KHC)
542 Z=Z+N-K
550 IF (.NOT. PTCO) GO TO 560
      Z=0
      KH0=IP*RL+KP+1
      DO 552 IS=1,KP
          K=IS-1
          KH=IKP+Z
          KH1=IP*FL+K+1
          A (KH) =A (KH) +P (KH1) *A (T+
552 Z=Z+N-K
      DO 554 J=KP,N
          KH=J+T-KP+1
554 A (KH) =A (KH) *P (KH0)
      Z=0
      DO 556 IS=1,IKP
          K=IS-1
          KH=IKP+Z
          A (KH) =A (KH) *P (KH0)
556 Z=Z+N-K
      RETURN
560 Z=0
      DO 562 IS=1,KP
          K=IS-1
          KH=KI+Z+1
          A (KH) =0. DO
562 Z=Z+N-K
      A (T+1) =1. DO /A (T+1)
      IF (IKP.GT.N) GO TO 564
      DO 568 J=IKP,N
          KH=J+T-KP+1
568 A (KH) =0. DO
564 RETURN
      END

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SUBROUTINE COLPIV
IMPLICIT REAL*8 (A-H,O-Z)
COMMON /OFUNCN/ A (100)
COMMON /OTFOS / LOS (50)
COMMON /KNUMBER/ KCO,KVA
COMMON /NUMTFC/ N,M
COMMON /MINIMO/ YG
COMMON /PIVOTT/ IPIV,KP
KP=LOS (1)
KHC=LOS (1) +1
YO=A (KH0)
IF (KVA.LT.2) RETURN
DO 602 K=2,KVA
    KH0=LOS (K) +1
    IF (A (KH0) .GT. YO) GO TO 602

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CUA04630
CUA04640
CUA04650
CUA04660
CUA04670
CUA04680
CUA04690
CUA04700
CUA04710
CUA04720
CUA04730
CUA04740
CUA04750
CUA04760
CUA04770
CUA04780
CUA04790
CUA04800
CUA04810
CUA04820
CUA04830
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CUA04870
CUA04880
CUA04890
CUA04900
CUA04910
CUA04920
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CUA04960
CUA04970
CUA04980
CUA04990
CUA05000
CUA05010
CUA05020
CUA05030
CUA05040
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CUA05060
CUA05070
CUA05080
CUA05090
CUA05100
CUA05110
CUA05120
CUA05130
CUA05140

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	YO=V (KH0)	CUA05150
	KP=LOS (K)	CUA05160
602	CONTINUE	CUA05170
	RETURN	CUA05180
	END	CUA05190
		CUA05200
	SUBROUTINE RATIO	CUA05210
	IMPLICIT REAL*8 (A-H,O-Z)	CUA05220
	INTEGER RL,V,Z	CUA05230
	COMMON /LENGTH/ PI	CUA05240
	COMMON /PIVOTE/ IP,KP	CUA05250
	COMMON /SUBJECT/ P (200)	CUA05260
	COMMON /OTPOS / KLOS (50) , LAS (50)	CUA05270
	COMMON /KNUMBER/ KCO	CUA05280
	COMMON /VALUES/ FPS ,MIO ,ITEP ,DSPLAY ,KIM	CUA05290
	COMMON /MINIMO/ YC	CUA05300
	COMMON /PARAME/ V	CUA05310
	COMMON /NUMERO/ N ,M	CUA05320
	IP=0	CUA05330
	IF (KCO.LT.1) RETURN	CUA05340
	DO 702 I=1,KCO	CUA05350
	KH=LAS (I) *PI+1	CUA05360
	KH1=KH+KP	CUA05370
	IF (V*P (KH1) .GT.EPS) GO TO 704	CUA05380
702	CONTINUE	CUA05390
	RETURN	CUA05400
704	YO=V*P (KH) /P (KH1)	CUA05410
	IP=LAS (I)	CUA05420
	Z=I+1	CUA05430
	IF (Z.GT.KCO) RETURN	CUA05440
	DO 706 I=Z,KCO	CUA05450
	KH=LAS (I) *PI+1	CUA05460
	KH1=KH+KP	CUA05470
	IF (V*P (KH1) .LE.EPS) GO TO 706	CUA05480
	YF=V*P (KH) /P (KH1)	CUA05490
	IF (YE.GE.YC) GO TO 708	CUA05500
	IP=LAS (I)	CUA05510
	YO=YE	CUA05520
	GO TO 706	CUA05530
708	IF (YF.NE.YC) GO TO 706	CUA05540
	IO=LAS (I)	CUA05550
	DO 710 K=1,N	CUA05560
	KH0=IP*PI+K+1	CUA05570
	KH2=IP*PI+KP+1	CUA05580
	KH=IO*PI+K+1	CUA05590
	YFP=V*P (KH0) /P (KH2)	CUA05600
	YEO=V*P (KH) /P (KH1)	CUA05610
	IF (YEP.LT.YFO) GO TO 706	CUA05620
	IF (YEO.LT.YFP) GO TO 712	CUA05630
710	CONTINUE	CUA05640
712	IP=IO	CUA05650
		CUA05660

706 CONTINUE  
RETURN  
END

SUBROUTINE CCONACT  
IMPLICIT REAL\*8 (A-H,O-Z)  
INTEGER PL,F  
COMMON /SUBJECT/ P(200)  
COMMON /PAFAME/ V,T,I0  
COMMON /NUMERC/ K1,I1  
COMMON /PIVOTE/ IP,KP  
COMMON /LFNGTH/ RL  
COMMON /FINAL / ISDONE  
KH=IP\*PL+KP+1  
IF(P(KH).EQ.0) GO TO 820  
PIVOTA=1. DO/P(KH)  
IIO=I0+1  
II1=I1+1  
KK0=1  
KK1=K1+1  
IF(II0.GT.II1) GO TO 802  
DO 804 II=IIO,II1  
I=II-1  
IF(I.EQ.IP) GO TO 804  
KH0=I\*PL+KP+1  
P(KH0)=P(KH0)\*PIVOTA  
DO 806 KK=KK0,KK1  
K=KK-1  
IF(K.EQ.KP) GO TO 806  
KH1=I\*PL+K+1  
KH2=IP\*PL+K+1  
P(KH1)=P(KH1)-P(KH2)\*P(KH0)  
806 CONTINUE  
804 CONTINUE  
DO 810 KK=KK0,KK1  
K=KK-1  
KH2=IP\*PL+K+1  
810 IF(K.NE.KP) P(KH2)=-P(KH2)\*PIVOTA  
802 P(KH)=PIVOTA  
RETURN  
820 ISDONE = 0  
RETURN  
END

SUBROUTINE SHOWIT  
IMPLICIT REAL\*8 (A-H,O-Z)  
DIMENSION SOL(50)  
INTEGER DSPLAY,PL,BASIC1,BASIC0  
COMMON /VALUES/ FPS,MIO,ITER,DSPLAY  
COMMON /NUMERIC/ N,M

CUA05670  
CUA05680  
CUA05690  
CUA05700  
CUA05710  
CUA05720  
CUA05730  
CUA05740  
CUA05750  
CUA05760  
CUA05770  
CUA05780  
CUA05790  
CUA05800  
CUA05810  
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CUA06010  
CUA06020  
CUA06030  
CUA06040  
CUA06050  
CUA06060  
CUA06070  
CUA06080  
CUA06090  
CUA06100  
CUA06110  
CUA06120  
CUA06130  
CUA06140  
CUA06150  
CUA06160  
CUA06170  
CUA06180

COMMON /OFUNCN/ A (100)	CUA06190
COMMON /SUBJCT/ P (200)	CUA06200
COMMON /BASICA/ BASIC0 (50) ,BASIC1 (50)	CUA06210
COMMON /LENGTH/ RL	CUA06220
IF (DSPLAY.EQ.0) GO TO 1890	CUA06230
IF (DSPLAY.EQ.1.AND.ITER.EQ.0) CALL MESSAG (17)	CUA06240
IF (DSPLAY.EQ.1.AND.ITER.EQ.0) GO TO 1820	CUA06250
IF (DSPLAY.GE.2.AND.ITER.EQ.0) GO TO 1812	CUA06260
1813 WRITE (6,213) ITER	CUA06270
IF (DSPLAY.EQ.1) GO TO 1820	CUA06280
1812 CALL PROFUN	CUA06290
IF (DSPLAY.EQ.2.AND.M.EQ.0.AND.ITER.EQ.0) GO TO 1820	CUA06300
CALL PRICON	CUA06310
1820 ITER = ITER + 1	CUA06320
WRITE (6,106) A (1)	CUA06330
K=M+1	CUA06340
DO 28 I=1,K	CUA06350
KIO=(I-1)*PL+1	CUA06360
SOL (I) =P (KIO)	CUA06370
28 CONTINUE	CUA06380
WRITE (6,107) (BASIC1 (K) ,SOL (K+1) ,K=1,M)	CUA06390
CALL MESSAG (18)	CUA06400
1890 RETURN	CUA06410
106 FORMAT (//,5X,' THE CURRENT VALUE OF THE OBJECTIVE FUNCTION IS ',	CUA06420
*F15.8)	CUA06430
213 FORMAT (//,15X,' ITERATION NUMBER ',I2)	CUA06440
107 FORMAT (//,5X,' THE VALUES OF THE VARIABLES ARE ',/,30X,'X',10 (I2,'	CUA06450
*=' ,F15.8,/,30X,'X'))	CUA06460
END	CUA06470
	CUA06480
	CUA06490
SUBROUTINE OPTIMC	CUA06500
IMPLICIT REAL*8 (A-H,O-Z)	CUA06510
DIMENSION SOL (50)	CUA06520
COMMON /VALUES/ FPS,MIO,ITER,DSPLAY	CUA06530
COMMON /FINAL / ISDONE	CUA06540
COMMON /NUM EPC/ N,M	CUA06550
COMMON /BASICA/ BASIC0 (50) ,BASIC1 (50)	CUA06560
COMMON /SUBJCT/ P (200)	CUA06570
COMMON /OFUNCN/ A (100)	CUA06580
COMMON /LENGTH/ RL	CUA06590
INTEGER DSPLAY,S,PL,BASIC1,BASIC0	CUA06600
IF (ISDONE.EQ.0) CALL MESSAG (19)	CUA06610
IF (ISDONE.EQ.0) GO TO 1750	CUA06620
IF (ISDONE.EQ.1) CALL MESSAG (20)	CUA06630
IF (DSPLAY.EQ.0.OR.DSPLAY.EQ.1) GO TO 1710	CUA06640
CALL PROFUN	CUA06650
CALL PRICON	CUA06660
1710 WRITE (6,106) A (1)	CUA06670
K=M+1	CUA06680
DO 28 I=1,K	CUA06690
KIO=(I-1)*PL+1	CUA06700

	SOL(I) = P(KIC)	CUA06710
28	CONTINUE	CUA06720
	WRITE(6,107) (BASIC1(K), SOL(K+1), K=1, M)	CUA06730
	CALL MESSAG(18)	CUA06740
	CALL MESSAG(21)	CUA06750
1750	RETURN	CUA06760
106	FORMAT(//,5X,' THE OPTIMAL IS ',F15.8)	CUA06770
107	FORMAT(/,5X,' THE VALUES OF THE VARIABLES ARE ',/,31X,'X',10(I2, 'CUA06780	
	*=',F15.8,/,31X,'X'))	CUA06790
108	FORMAT(/,5X,' ALL THE OTHERS ARE EQUAL TO ZERO ',/,78('-'),/,58X,'CUA06800	
	10SCAF MANRIQUE, 1982',/,78('-'))	CUA06810
	END	CUA06820
		CUA06830
		CUA06840
	SUBFOUNTINE DIALOG	CUA06850
	IMPLICIT REAL*8 (A-H,O-Z)	CUA06860
	DIMENSION G(200), D(100)	CUA06870
	INTEGER DISPLAY	CUA06880
	COMMON /NUMREQ/ N, M	CUA06890
	COMMON /SUBJECT/ P(200), Q(50)	CUA06900
	COMMON /OPUNCN/ A(100), B(50), C	CUA06910
	COMMON /VALUES/ F, MIO, ITER, DISPLAY	CUA06920
	DATA BLANK/' '/	CUA06930
	CALL MESSAG(1)	CUA06940
	CALL MESSAG(2)	CUA06950
	CALL MESSAG(3)	CUA06960
	KEY=0	CUA06970
100	CALL MESSAG(4)	CUA06980
	CALL READIT(N)	CUA06990
	IF(N.GT.0) GO TO 41	CUA07000
	WRITE(6,60) N	CUA07010
	GO TO 100	CUA07020
41	IF(KEY.EQ.1) GO TO 1901	CUA07030
102	CALL MESSAG(5)	CUA07040
	CALL READIT(M)	CUA07050
	IF(M.GE.0) GO TO 42	CUA07060
	CALL MESSAG(6)	CUA07070
	GO TO 102	CUA07080
42	IF(M.LE.N) GO TO 43	CUA07090
	WRITE(6,67) N	CUA07100
	GO TO 102	CUA07110
43	IF(KEY.EQ.1) GO TO 196	CUA07120
	CALL MESSAG(7)	CUA07130
	KO=0	CUA07140
	JIK=0	CUA07150
121	DO 13 J=1, N	CUA07160
	IF(JIK.EQ.N) GO TO 104	CUA07170
	IF(JIK.NE.0) GO TO 122	CUA07180
	JL=J-1	CUA07190
	NI=N-JL	CUA07200
	NI1=NI	CUA07210
	IF(NI.EQ.1) GO TO 104	CUA07220

	WRITE(6,70)NI,J	CUA07230
	GO TO 105	CUA07240
122	WRITE(6,70)NIL,JIK	CUA07250
	GO TO 105	CUA07260
104	CALL MESSAG(8)	CUA07270
105	CALL PEADEM(D,NIL)	CUA07280
	KIN=1	CUA07290
	IF(JIK.NE.C)KC=IO-NIL-1	CUA07300
	IF(JIK.NE.C)NI=NIL	CUA07310
	DO 14 K=KIN,NI	CUA07320
	KO=KO+1	CUA07330
	A(KO)=D(K)	CUA07340
14	CONTINUE	CUA07350
	IF(KEY.EQ.1)IA=IA-NIL	CUA07360
	IF(KEY.FQ.1)GO TO 1017	CUA07370
13	CONTINUE	CUA07380
118	WRITE(6,68)N	CUA07390
	CALL PEADEM(B,N)	CUA07400
	IF(KEY.EQ.1)GO TO 115	CUA07410
116	CALL MESSAG(9)	CUA07420
	CALL PEADEM(C,1)	CUA07430
	IF(KEY.EQ.1)GO TO 113	CUA07440
	IF(M.EQ.0)GO TO 140	CUA07450
130	JA = 1	CUA07460
	JIO = 0	CUA07470
112	DO 11 I=1,M	CUA07480
	IF(JTO.NE.C.AND.KEY.EQ.1)GO TO 110	CUA07490
	WRITE(6,66)N,I	CUA07500
	GO TO 111	CUA07510
110	WRITE(6,66)N,JIO	CUA07520
111	CALL PEADEM(G,N)	CUA07530
	IF(KEY.FQ.0)JIO=I	CUA07540
	IF(KEY.FQ.1)JA=(JTO-1)*N+1	CUA07550
	JJ=JIO*N	CUA07560
	J=1	CUA07570
	DO 12 K=JA,JJ	CUA07580
	D(K) = G(J)	CUA07590
	J = J + 1	CUA07600
12	CONTINUE	CUA07610
	IF(KEY.FQ.1)GO TO 108	CUA07620
	JA = JJ + 1	CUA07630
11	CONTINUE	CUA07640
	IF(M.GT.1)GO TO 103	CUA07650
	CALL MESSAG(10)	CUA07660
	GO TO 131	CUA07670
103	WRITE(6,64)M	CUA07680
131	CALL PEADEM(Q,M)	CUA07690
	DO 10 I=1,M	CUA07700
	IF(Q(I).GT.0)GO TO 10	CUA07710
	WRITE(6,65)I,Q(I)	CUA07720
	GO TO 103	CUA07730
10	CONTINUE	CUA07740

	IF (KEY.EQ.1) GO TO 180	CUA07750
140	CALL MESSAG(22)	CUA07760
	CALL BFADEM(E,1)	CUA07770
	IF (KEY.EQ.1) GO TO 181	CUA07780
	----- NEXT SECTION CHECKS THE CORRECTNESS OF THE DATA -----	CUA07790
	CALL MESSAG(11)	CUA07800
	KFY = 1	CUA07810
1901	WRITE(6,901) N	CUA07820
	REWIND 5	CUA07830
	READ(5,33,END=106) CHAR	CUA07840
	IF (CHAR.NE.BLANK) GO TO 100	CUA07850
106	IF (M.EQ.0) CALL MESSAG(12)	CUA07860
	IF (M.EQ.0) GO TO 701	CUA07870
	WRITE(6,74) M	CUA07880
701	REWIND 5	CUA07890
	READ(5,33,END=107) CHAR	CUA07900
	IF (CHAR.NE.BLANK) GO TO 102	CUA07910
107	IC=1	CUA07920
	IA=0	CUA07930
	DO 16 I=1,N	CUA07940
	IL=I-1	CUA07950
	MII=N-IL	CUA07960
	IO=IO+NIL	CUA07970
1017	IF (I.EQ.N) GO TO 128	CUA07980
	IA=IA+NIL	CUA07990
	JO=IO-NIL	CUA08000
	WRITE(6,79) NIL,I,(A(KI),KI=JO,IA)	CUA08010
	GO TO 119	CUA08020
128	LST=(N+1)*N/2	CUA08030
	WRITE(6,80) A(LST)	CUA08040
119	REWIND 5	CUA08050
	READ(5,33,END=16) CHAR	CUA08060
	IF (CHAR.EQ.BLANK) GO TO 16	CUA08070
	JIK=I	CUA08080
	GO TO 121	CUA08090
16	CONTINUE	CUA08100
115	WRITE(6,78) N,(B(KI),KI=1,N)	CUA08110
	REWIND 5	CUA08120
	READ(5,33,FND=113) CHAR	CUA08130
	IF (CHAR.NE.BLANK) GO TO 118	CUA08140
113	WRITE(6,77) C	CUA08150
	REWIND 5	CUA08160
	READ(5,33,FND=1180) CHAR	CUA08170
	IF (CHAR.NE.BLANK) GO TO 116	CUA08180
1180	IF (M.EQ.0) GO TO 181	CUA08190
	DO 15 IK=1,M	CUA08200
	KA=IK*N	CUA08210
	KO=(IK-1)*N+1	CUA08220
108	WRITE(6,75) N,IK,(P(J),J=KO,KN)	CUA08230
	REWIND 5	CUA08240
	READ(5,33,FND=15) CHAR	CUA08250
	IF (CHAR.EQ.BLANK) GO TO 15	CUA08260

	JIC=IK	CUA08270
	GO TO 112	CUA08280
15	CONTINUE	CUA08290
180	WRITE(6,72) M, (Q(I), I=1, M)	CUA08300
	REWIND 5	CUA08310
	FEAD(5,33,END=181) CHAF	CUA08320
	IF(CHAF.NE.BLANK) GO TO 103	CUA08330
181	WRITE(6,81) F	CUA08340
	REWIND 5	CUA08350
	FEAD(5,33,END=888) CHAF	CUA08360
	IF(CHAF.NE.BLANK) GO TO 140	CUA08370
888	CALL MESSAG(13)	CUA08380
	CALL FEADIT(DSPLAY)	CUA08390
	RETURN	CUA08400
60	FORMAT(/, ' YOUR VALUE OF N = ', I2, ' IS NOT POSITIVE ! ', /, ' )	CUA08410
	1PSTAMP .. ')	CUA08420
67	FORMAT(/, ' THE STANDARDIZED PROBLEM REQUIRES M <= N ... ', /, ' )	CUA08430
	* ' IN YOUR PROBLEM, A MAXIMUM OF ', I2, ' CONSTRAINTS (M) WILL BE ACCCUA08440	
	*EPTD')	CUA08450
68	FORMAT(/, ' TYPE THE ', I2, '-VECTOR B, (THE LINEAR PART OF YOUR OBJECTUA08460	
	1CTIVE FUNCTION) ' )	CUA08470
66	FORMAT(/, ' TYPE THE ', I2, ' VALUES OF THE ROW NO. ', I2, ' IN YOUR COCUA08480	
	1NSTRAINTS MATRIX P')	CUA08490
70	FORMAT(/, ' TYPE THE ', I2, ' VALUES OF THE ROW NUMBER ', I2)	CUA08500
64	FORMAT(/, ' TYPE THE VALUE OF YOUR ', I2, '-VECTOR Q. (THE RIGHT-HANDCUA08510	
	1 SIDE VALUE IN YOUR ', /, ' CONSTRAINTS) ' )	CUA08520
65	FORMAT(/, ' THE COMPONENT NO. ', I2, ' OF YOUR Q VECTOR, NAMELY ', D12CUA08530	
	* .5, ' IS NOT POSITIVE: ', /, ' YOUR PROBLEM IS NOT IN STANDARD FORM', /CUA08540	
	* , ' INPUT AGAIN .... ' )	CUA08550
79	FORMAT(/, ' ARE THESE THE CORRECT ', I2, ' ELEMENTS OF THE ROW NUMBERSCUA08560	
	* ', I2, ' OF MATRIX A ', /, ' IN YOUR OBJECTIVE FUNCTION ? ', /, 1X, 99D1CUA08570	
	* 2.5)	CUA08580
74	FORMAT(/, ' YOU HAVE ', I2, ' CONSTRAINTS ? ' )	CUA08590
33	FORMAT(A5)	CUA08600
901	FORMAT(/, ' YOU HAVE ', I2, ' VARIABLES ? ' )	CUA08610
78	FORMAT(/, ' ARE THESE THE CORRECT VALUES OF YOUR ', I2, '-VECTOR B ', CUA08620	
	* /, ' (THE LINEAR PART IN YOUR OBJECTIVE FUNCTION) ? ', /, 50 (2X, D12.5) )CUA08630	
80	FORMAT(/, ' IS ', D12.5, ' THE LAST ELEMENT OF THE MATRIX IN YOUR OBJCUA08640	
	* ECTIVE FUNCTION ? ' )	CUA08650
72	FORMAT(/, ' ARE THESE THE CORRECT VALUES OF YOUR ', I2, '-VECTOR Q ' CUA08660	
	* /, ' (THE RIGHT-HAND SIDE TERMS IN YOUR CONSTRAINTS) ? ', /, CUA08670	
	* 10 (2X, D12.5) )	CUA08680
75	FORMAT(/, ' ARE THESE THE CORRECT ', I2, ' VALUES OF ROW NUMBER ', I2, CUA08690	
	1 /, ' IN YOUR MATRIX P (CONSTRAINTS) ? ', /, 99 (2X, D12.5) )	CUA08700
77	FORMAT(/, ' THE CONSTANT VALUE C IN YOUR OBJECTIVE FUNCTION IS ' CUA08710	
	1, D12.5, ' ? ' )	CUA08720
81	FORMAT(/, ' THE WIDTH OF YOUR ZERO-BAND IS ', D12.5, ' ' ? ' )	CUA08730
	END	CUA08740
		CUA08750
		CUA08760
	SUBROUTINE MESSAG(INDEX)	CUA08770
	GO TO (1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 21, 22) ,	CUA08780

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*INDFX
1 WRITE(6,90)
  RETURN
2 WRITE(6,92)
  RETURN
3 WRITE(6,93)
  RETURN
4 WRITE(6,101)
  RETURN
5 WRITE(6,61)
  RETURN
6 WRITE(6,62)
  RETURN
7 WRITE(6,69)
  RETURN
8 WRITE(6,71)
  RETURN
9 WRITE(6,67)
  RETURN
10 WRITE(6,59)
  RETURN
11 WRITE(6,73)
  RETURN
12 WRITE(6,70)
  RETURN
13 WRITE(6,88)
  RETURN
14 WRITE(6,63)
  RETURN
15 WRITE(6,212)
  RETURN
16 WRITE(6,302)
  RETURN
17 WRITE(6,181)
  RETURN
18 WRITE(6,108)
  RETURN
19 WRITE(6,105)
  RETURN
20 WRITE(6,140)
  RETURN
21 WRITE(6,141)
  RETURN
22 WRITE(6,150)
  RETURN
90 FOPMAT(//,59X,'UNIVERSITY OF OTTAWA',//,62X,'SYSTEMS SCIENCE',//,69XCUA09240
*, '1982',//,1X,78(' - '),//,35X,' CUADPC',//,1X,78(' - '),//, ' THIS PACHA09250
*CKAGE USES THE METHOD OF BEALE TO SOLVE THE QUADRATIC PROGRAMMING'CUA09260
* ,//, ' PROBLEM, HAVING N VARIABLES AND M CONSTRAINTS (M<=N), GIVEN TCUA09270
* N ITS STANDARD',//, ' FORM :',//, CU09280
* ' MIN G = (1/2) X'AX + B'Y + C',//, CUA09290
* ' SUBJECT TO : PX <= Q , WITH Q(J) > 0',//, CUB09300

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*'           X(K) >= 0                               K = 1, 2, ... , N ',/, CUA09310
*'           J = 1, 2, ... , M ',/, CUA09320
*' WHEREF :   C : THE CONSTANT VALUE OF YOUR OBJECTIVE FUNCTION',/, CUA09330
*'           B : THE VECTOR OF THE LINEAR PART IN THE O. F. ',/, CUA09340
*'           A : AN NXN MATRIX CONTAINING THE QUADRATIC COEFFICIENTS',/, CUA09350
* OF THE O. F.',/, CUA09360
*'           P : AN NXM MATRIX WITH THE COEFFICIENTS OF THE CONSTRAINTS',/, CUA09370
*NTS ',/, CUA09380
*'           Q : AN M-VECTOR WITH THE RIGHT-HAND SIDE VALUES OF THE CONSTRAINTS',/, CUA09390
* CONSTRAINTS',/, CUA09400
92 FORMAT(' NOTE THAT MATRIX A IS A SYMMETRIC MATRIX WITH A(I,J) = ', CUA09410
*,/, CUA09420
* ' TWICE THE COEFFICIENT OF X(K)**2 WHEN K=I=J (THE DIAGONAL OF THE', CUA09430
* MATRIX)',/, CUA09440
* ' COEF. OF X(I)X(J) WHEN I DIFFERENT FROM J (THE OFF-DIAGONAL ELC', CUA09450
* ELEMENTS) ') CUA09460
93 FORMAT(/, ' WHEN TYPING THE DATA REQUESTED, PLEASE RECALL : ',/, CUA09470
* ' * AT LEAST ONE BLANK (SPACE) MUST SEPARATE SUCCESSIVE NUMBERS', CUA09480
* ',/, CUA09490
* ' * THE DECIMAL POINT IS NOT NECESSARY, UNLESS IT IS PART OF THE', CUA09500
* NUMBER ',/, CUA09510
* ' * IF YOU WANT TO QUIT TYPE "Q" INSTEAD OF THE INFORMATION REQUESTED', CUA09520
* ',/, CUA09530
* ' * WHEN ANSWERING TRUE OR FALSE, ',/, CUA09540
* ' TYPE "NO" AND <CR> IF FALSE, OR ',/, CUA09550
* ' <CR> (HIT RETURN) IF TRUE ',/, CUA09560
101 FORMAT(/, ' TYPE THE VALUE OF N (I.E. THE NUMBER OF VARIABLES YOU', CUA09570
* HAVE) ') CUA09580
61 FORMAT(/, ' TYPE THE VALUE OF M (I.E. THE NUMBER OF CONSTRAINTS YOU', CUA09590
* HAVE) ') CUA09600
62 FORMAT(/, ' YOUR NUMBER OF CONSTRAINTS CANNOT BE NEGATIVE !',/, CUA09610
* TRY AGAIN .. ') CUA09620
69 FORMAT(/, ' THE QUADRATIC PART OF THE OBJECTIVE FUNCTION, BEING A', CUA09630
* SYMMETRIC MATRIX, ',/, ' ONLY THE UPPER-TRIANGULAR PART MUST BE TYPED', CUA09640
* ') CUA09650
71 FORMAT(/, ' TYPE THE VERY LAST ELEMENT OF YOUR MATRIX ') CUA09660
67 FORMAT(/, ' TYPE THE VALUE OF C, I.E. THE CONSTANT PART IN YOUR OBJECTIVE', CUA09670
* FUNCTION ') CUA09680
59 FORMAT(/, ' TYPE THE VALUE OF Q IN YOUR CONSTRAINT (I.E., THE RIGHT-HAND', CUA09690
* SIDE VALUE) ') CUA09700
73 FORMAT(/, ' NOW, LET US CHECK THE CORRECTNESS OF THE DATA RECEIVED', CUA09710
*,/, ' TYPE "NO" AND <CR> IF FALSE, OR ', CUA09720
*,/, ' <CR> (HIT RETURN) IF TRUE ') CUA09730
700 FORMAT(/, ' YOUR PROBLEM IS UNCONSTRAINED ?') CUA09740
38 FORMAT(/, ' SELECT ONE OF THE FOLLOWING OPTIONS : ',/, ' 0 >ONLY THE', CUA09750
* FINAL ANSWER WILL BE SHOWN',/, ' 1 >THE VALUES OF THE', CUA09760
* OBJECTIVE FUNCTION ', CUA09770
* ' AND THE BASIC VARIABLES WILL BE SHOWN ',/, ' 5X, AT EACH ITERATION', CUA09780
* ',/, ' 2 >THE COMPLETE TABLEAU FOR THE METHOD OF BRUTE WILL BE', CUA09790
* DISPLAYED AT EACH ',/, ' 5X, ITERATION ',/, ' TYPE THE NUMBER OF THE', CUA09800
* OPTION YOU CHOOSE') CUA09810
63 FORMAT(' YOU DID NOT ENTER ANY DATA ...! INPUT AGAIN OR TYPE Q TO', CUA09820

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* QUIT')
212 FORMAT(/,15X,'ORIGINAL TABLEAU')
302 FORMAT(//,5X,'CONSTRAINTS')
181 FORMAT(/,15X,'STARTING POINT ')
108 FORMAT(/,5X,' ALL THE OTHERS ARE EQUAL TO ZERO ',/,30(' '))
105 FORMAT(//,5X,' *** THE PROBLEM DOES NOT HAVE A FINITE SOLUTION **')
** ',/,78(' '),/,58X,'OSCAF MANRIQUE, 1982',/,78(' '))
140 FORMAT(///,36X,' SOLUTION ')
141 FORMAT(58X,'OSCAF MANRIQUE,1982',/,79(' '))
150 FORMAT(/,' THE VALUES IN THE ZERO-BAND ARE TAKEN AS ZERO IN CRITICCUA
*AL POINTS ',/, ' OF THE PROBLEM.',/, ' TYPE THE VALUE OF THE ZERO-BAND
*ND FOR YOUR PROBLEM ')
END

SUBROUTINE READFM(F,L)
INTEGER A(80),NUMER(12),DOT,BLANK,QUIT
INTEGER FMT(40)
DOUBLE PRECISION F(40)
LOGICAL FLAG
DATA BLANK/' ',A/80*' ',/,QUIT/'Q',/,MIN/'-'/,DOT/'.'/,
DATA NUMER/'1','2','3','4','5','6','7','8','9','10','11','12'/
40 REWIND 5
READ(5,100,END=77) (A(I),I=1,80)
FLAG=.TRUE.
DO 45 I=1,80
IF(A(I).EQ.QUIT)CALL EXIT
45 CONTINUE
77 DO 60 J = 1,80
IF(A(J).NE.BLANK)GO TO 63
60 CONTINUE
CALL MESSAGE(14)
GO TO 40
63 DO 70 I=1,80
IF(A(I).EQ.MIN.OR.A(I).EQ.DOT.OR.A(I).EQ.BLANK)GO TO 70
DO 72 K=1,12
IF(A(I).EQ.NUMER(K))GO TO 70
72 CONTINUE
WRITE(6,101) A(I)
FLAG=.FALSE.
GO TO 75
70 CONTINUE
75 IF(FLAG)GO TO 66
WRITE(6,102) (A(K),K=1,80)
GO TO 40
66 JAK=1
I2=0
I1 = 0
K=1
50 IF(A(K).EQ.BLANK.AND.I1.NE.0)I2=K-1
IF(A(K).NE.BLANK.AND.I1.EQ.0)I1=K
IF(I2.EQ.0)GO TO 10

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CPA09830
CPA09840
CPA09850
CPA09860
CPA09870
CPA09880
CPA09890
CPA09900
CPA09910
CPA09920
CPA09930
CPA09940
CPA09950
CPA09960
CPA09970
CPA09980
CPA09990
CPA10000
CPA10010
CPA10020
CPA10030
CPA10040
CPA10050
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CPA10180
CPA10190
CPA10200
CPA10210
CPA10220
CPA10230
CPA10240
CPA10250
CPA10260
CPA10270
CPA10280
CPA10290
CPA10300
CPA10310
CPA10320
CPA10330
CPA10340

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GC TO 75	CMA 10870
70 CONTINUE	CMA 10880
75 IF (FLAG) GO TO 66	CMA 10890
WRITE (6, 102) (A(K), K=1, 80)	CUA 10900
GO TO 40	CUA 10910
66 MIKE=0	CMA 10920
I2=0	CUA 10930
I1=0	CUA 10940
K=1	CUA 10950
50 IF (A(K).EQ.BLANK.AND.I1.NE.0) I2=K-1	CUA 10960
IF (A(K).NE.BLANK.AND.I1.EQ.0) I1=K	CMA 10970
IF (I2.EQ.0) GO TO 10	CUA 10980
WRITE (99, 101) (A(KI), KI=I1, I2)	CUA 10990
CALL FFCB99 (FMT, 10)	CUA 11000
IH=I2-I1+1	CMA 11010
WRITE (99, 99) IH	CUA 11020
CALL FFCB99	CMA 11030
READ (99, FMT) F	CUA 11040
MIKE=1	CUA 11050
10 K=K+1	CMA 11060
IF (K.GT.80) GO TO 78	CUA 11070
IF (MIKE.NE.0) GO TO 78	CUA 11080
GO TO 50	CUA 11090
78 RETURN	CMA 11100
99 FORMAT (' (I', I1, ')')	CUA 11110
101 FORMAT (9A1)	CUA 11120
100 FORMAT (80A1)	CUA 11130
102 FORMAT (' ***** BEFCE IN YOUR ENTRY :', /, 80A1, /, ' INPUT AGAIN OR', *TYPE Q TO QUIT')	CUA 11140
103 FORMAT (' -----> ', A1)	CUA 11150
END	CUA 11160
	CMA 11170

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